

Aim: Get large cardinal strength from  $\omega$  All cardinals  $> \omega$  are singular  $\aleph_\alpha$  (concluded by work of Moti Gitik)

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We will use:

Core model thy      descriptive set thy

core model induction +  $\varphi \rightsquigarrow$  lower bounds

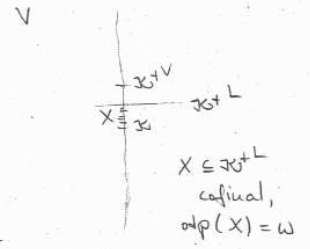
(why not just core model thy?)

Why does  $\varphi$  give  $O^*$ ? (J. covering Lemma)

No inner model with choice can compute  $\aleph^+$  correctly, for all  $\aleph$ .

$L \subset \text{HOD}[X]$ .

Jensen's Cov. Lemma yields Either  $\exists Y \geq X, \bar{Y} \leq \aleph, Y \in L$  or else  $O^* \text{ ex}$ .



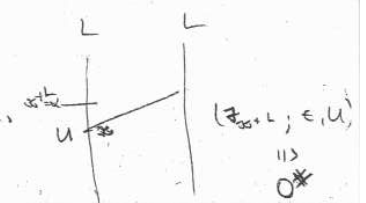
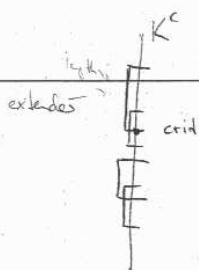
Properties of  $L$  used in this proof:

- $L^{\text{HOD}} = L^{\text{HOD}[X]}$  where  $X$  is gen. / HOD forcing absoluteness of  $L$
- weak covering for  $L$ , i.e., if  $\aleph$  is singular cardinal, then  $\aleph^+ \leq \aleph^+$  (?)

Now we want to produce a model  $K$  which has these properties, but which also has  $O^*$  in it, etc.

How do we produce such a  $K$ ?

First produce  $K^c$ , a preliminary version of  $K$ .  $c$  stands for "certified".

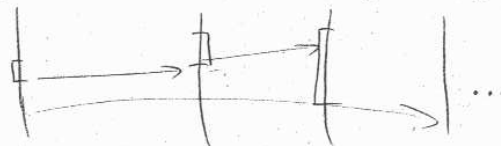


$K^c$  is not forcing absolute, also  $K^c$  does not seem to satisfy weak covy (in general)

Using  $K^c$ , we want to take a hull  $K \cong X \subset K^c$ .

In order to do that, we need the full iterability of  $K^c$ , we need:

$U_{\text{crit}_0} = K^c$        $U_{\text{crit}_1}$        $U_{\text{crit}_2}$        $U_{\text{crit}_3}$



yields tree structure



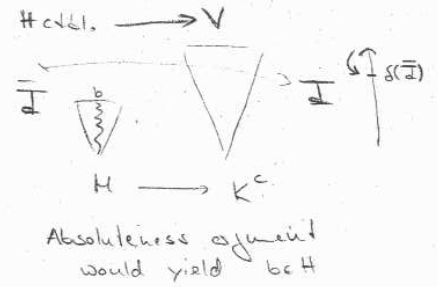
one can always pick the right branch, if one did not do a mistake before

$K^c$  is "weakly iterable": if  $M \rightarrow K^c$ ,  $M$  countable, then there is an iteration strategy for  $M$  w.r.t. countable iteration trees

We want:  $K^c$  is fully iterable

If  $H^V$  is closed under  $\mathcal{Q}$ -structures, then (1)  $\Rightarrow$  (2).

$\mathcal{Q} \leq M_b^I$   
 $\mathcal{Q} \neq S(I)$  is not woodin

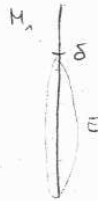


However, (\*) is sometimes false, and in fact  $K^c$  need not be fully iterable.

Look at  $M_1$ :

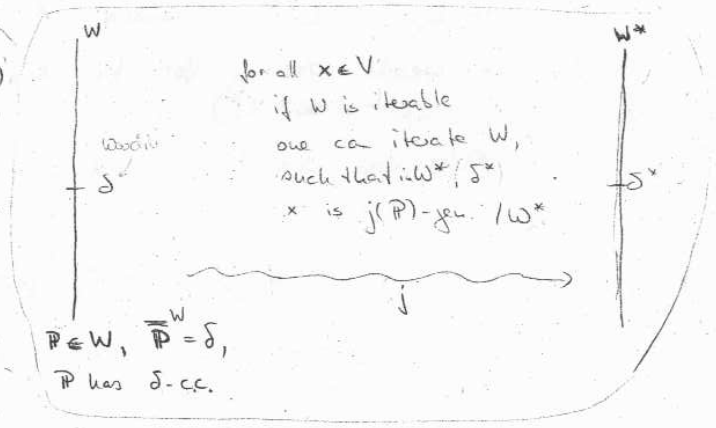
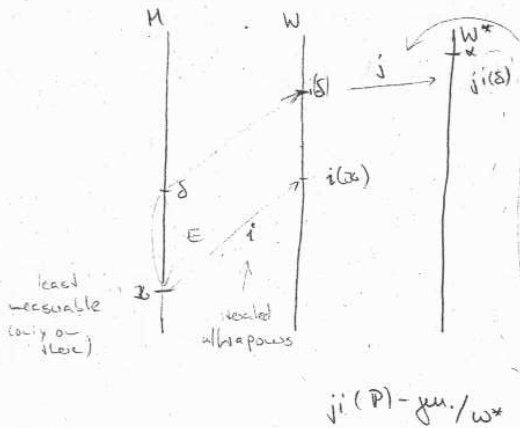
$M_1 \models \delta$  is Woodin

$L[E]$ , where  $E \in \delta$



We prove, that  $M_1$  (is not iterable) does not know, how to iterate itself:

Suppose  $M_1 \models$  "I'm fully iterable".



$W^*[E] = L[E]$

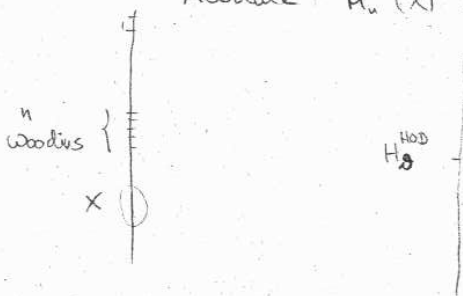
$\alpha = j(i(\delta^{M_1})) = j(i(\delta)^{W^*[E]}) = j(i(\delta)^{M_1})$   
 $\Rightarrow j'' \delta^{M_1}$  is cof in  $j(i(\delta)^{M_1})$   
 (nonsense for large enough  $\delta$ )

Core model induction:

we showed:  $\varphi \Rightarrow O^*$  ex. in fact  $X^*$  ex for all  $X$

First  $w$  induction steps: Use  $\varphi$  to show:

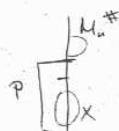
Assume  $M_n^*(X)$  ex. f.a.  $X$ . Show  $M_{n+1}^*(X)$  ex. f.a.  $X$ .



$L^{M_n^*}(H_\delta^{HOD})$  Look at  $K^c(X) \cap L^{M_n^*}(H_\delta^{HOD}) =: W$

Possibilities:

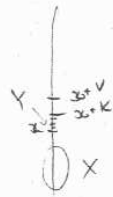
$\bullet$   $W$  has a Woodin cardinal



so we're finished

•  $W$  is fully iterable

isolate (using full iterability)  $K(X) \prec^{M^*} (H_\beta^{HOD})$



$L^{M^*}(H_\beta^{HOD}) [Y] \dots (?)$

Construct  $K$  + use  $\varphi$  to get  $\bar{Y}$

So  $\forall u \forall X \ M^*(X) \text{ ex.} \Rightarrow \text{Proj. Det.} = \text{PD}$

$\text{PD} = J_2(\mathbb{R}) \models \text{AD}$

$J_1(\mathbb{R}) = V_{\omega+1}$   
 $J_2(\mathbb{R}) = \text{red. cl. of } J_1(\mathbb{R}) \cup J_1(\mathbb{R})$

Now we want to prove inductively

$J_\alpha(\mathbb{R}) \models \text{AD}$  f.o.  $\alpha$ , using  $\varphi$ .

So  $J_\beta(\mathbb{R}) \cap J_\alpha(\mathbb{R}) = \text{the proj. sets}$

If  $\alpha$  is the least one with

$J_\alpha(\mathbb{R}) \models \neg \text{AD}, \alpha = \beta + 1, \text{ some } \beta$

In " $\neg \text{AD}$ " is  $\Sigma_1$ , so  $J_\alpha(\mathbb{R}) \models$  a  $\Sigma_1$ -statement which is false in all  $J_\beta(\mathbb{R}), \beta < \alpha$

Def.  $[\alpha, \beta]$  is a  $\Sigma_1$ -gap iff  $J_\alpha(\mathbb{R}) \prec_{\Sigma_1}^R J_\beta(\mathbb{R})$ , and

there is no  $\bar{\alpha} < \alpha$  s.t.

$J_{\bar{\alpha}}(\mathbb{R}) \prec_{\Sigma_1}^R J_\alpha(\mathbb{R})$ .

There's no  $\bar{\beta} > \beta$

$J_\beta(\mathbb{R}) \prec_{\Sigma_1}^R J_{\bar{\beta}}(\mathbb{R})$ .

The gaps partition the ordinals.

The first place where AD fails begins a gap.

THEOREM (Kechris-Woodin).  $\text{AD}^{L(\mathbb{R})}$  follows from:  $L(\mathbb{R}) \models$  every set which has scale is determined

DEF.  $\alpha$  is critical iff there is a set  $A \in \mathbb{R}$  which has scales in  $J_{\alpha+1}(\mathbb{R})$  but not in  $J_\alpha(\mathbb{R})$ .

By the K.-W.-Thm., it suff. to show: If  $\alpha$  is critical, then  $J_{\alpha+1}(\mathbb{R}) \models \text{AD}$

Critical ordinals:

use certain nice, obtain structures closed under  $\mathcal{Q}$ -operator

First Case:  $\alpha$  is  $\mathbb{R}$ -inadmissible and

- (a)  $\text{cf}(\alpha) = \omega_1$ ,
- (b)  $\text{cf}(\alpha) = \omega$  or
- (c)  $\alpha$  is successor ordinal, but the previous gap is not strong

Second Case: "end of the gap Case"

- (a)  $\alpha$  ends a proper weak gap,  $[\bar{\alpha}, \alpha]$  or
- (b)  $\alpha$  is a succ. ordinal and the previous gap is strong

Thm. (Deemman) .. Let  $W$  be an inner model,  $W \models \delta$  is Woodin,  
 $\text{Col}(\delta) \cap W = W$ . Let  $A \in \mathbb{R}$ . Supp.  $\exists \mathcal{T} \in W^{\text{Col}(W, \delta)}$  and  $\exists$  it. stat.  $\Sigma$   
 for  $W$  s.t. whenever  $i: W \rightarrow W^*$  is according to  $\Sigma$ , then for all  
 $g \in \text{Col}(W, i(\delta))$ -gen. /  $W^*$ ,  $A \cap W^*[g] = i(\mathcal{T})^g$ .

Then  $A$  is determined.

The argument in Case 2(a):

Starting point: We have a ctbl.  $N$

for  
 $A \in \mathbb{R}^N$ ,  
 there is  
 $\mathcal{T} \in N^{\text{Col}(W, \delta)}$   
 "capturing"  $A$



} closed under  
 $C_{\Sigma}^{\mathbb{R}^N}$

$[\bar{x}, x]$

$$g_u(\mathbb{R}^N) = \mathbb{R}$$

$$B \in \Sigma_u^{\mathbb{R}^N} \Leftrightarrow B = \bigcup_{u \in W} A_u \text{ w.t. } A_u \in \mathbb{R}^N$$



Go for a mouse  $M$  s.t.

- $U \in M$ ,  $M \models U$  ctbl.,  $M$  can see  $(\mathcal{T}_{A_u} \mid u \in W)$
- $M \models \delta$  is Woodin
- $M$  can see  $\Sigma \cap M$ , as given by  $\dot{\Sigma}^M = \Sigma \cap M$
- in fact  $M$  is iterable via  $\Gamma$  s.t. if  
 $j: M \rightarrow M^*$  is acc. to  $\Gamma$ , then  $j(\dot{\Sigma})^{M^*} = \Sigma \cap M^*$
- in fact, if  $j: M \rightarrow M^*$  is (like) above and  $g$  is  
 $\text{Col}(W, j(\delta))$ -gen. /  $M^*$ , then from  $j(\dot{\Sigma})^{M^*}$  you can  
 "easily" read off  $\Sigma \cap M^*[g]$ .

Claim:  $M$  is a  $\tau$ -mouse capturing  $B = \bigcup A_u$ .

( "hybrid mouse" )