What is...Hadwiger's conjecture?

Or: Coloring is difficult...

## Four colors suffice



- ▶ Four color theorem (4CT) Any map can be colored with four colors
- Proposes in 1852 when Guthrie tried to color the map of counties of England, proved in 1976 by Appel–Haken
- ▶ The proof is one of the main achievements in graph theory

## A different perspective on the 4CT



Complete bipartite graph on 6 vertices  $K_{3,3}$  Complete graph on 5 vertices  $K_5$ 

- Wagner's theorem A graph is planar  $\Leftrightarrow$  no  $K_{3,3}$  or  $K_5$  minors
- Coloring = adjacent vertices get different colors
- 4CT reformulated(?) A graph with at most a  $K_4$  minor is 4-colorable

Where is  $K_{3,3}$ ?



• Not that  $K_{3,3}$  is bipartite = 2-colorable

- Any graph with  $\geq$  one edge needs at least 2 colors
- ▶ Hence,  $K_{3,3}$  should not play any role for the colorability

Conjecture (1943)

If G is loopless and has no  $K_t$  minor then its chromatic number is < t

- Theorem The case t = 5 is true (4CT proven 1976)
- Theorem The case t = 6 is true (proven in 1993)
- ▶ t > 6 is open, but: Theorem The conjecture is almost always true:

## Hadwiger's Conjecture is True for Almost Every Graph

## B. BOLLOBÁS, P. A. CATLIN\* AND P. ERDÖS

The contraction clique number ccl(G) of a graph G is the maximal r for which G has a subcontraction to the complete graph  $K^r$ . We prove that for d > 2, almost every graph of order n satisfies  $n((\log_2 n)^{\frac{1}{2}}+4)^{-1} \le ccl(G) \le n(\log_2 n - d \log_2 \log_2 n)^{-\frac{1}{2}}$ . This inequality implies the statement in the title.

The bound need not to be sharp



• The Petersen graph contains a  $K_5$  minor

► The Petersen graph is 3 -colorable

Thank you for your attention!

I hope that was of some help.