## What is...Hadwiger's conjecture?

Or: Coloring is difficult...



Four color theorem (4CT) Any map can be colored with four colors

- Proposes in 1852 when Guthrie tried to color the map of counties of England, proved in 1976 by Appel-Haken
- The proof is one of the main achievements in graph theory


## A different perspective on the 4CT

Complete bipartite graph on 6 vertices $K_{3,3}$ Complete graph on 5 vertices $K_{5}$


Wagner - top to bottom A graph is planar if and only if it does not contain

$$
K_{3,3} \text { or } K_{5} \text { as a minor }
$$




Wagner's theorem A graph is planar $\Leftrightarrow$ no $K_{3,3}$ or $K_{5}$ minors

- Coloring $=$ adjacent vertices get different colors
- 4CT reformulated(?) A graph with at most a $K_{4}$ minor is 4 -colorable


## Where is $K_{3,3}$ ?



- Not that $K_{3,3}$ is bipartite $=2$-colorable
- Any graph with $\geq$ one edge needs at least 2 colors
- Hence, $K_{3,3}$ should not play any role for the colorability


## Enter, the theorem(s)

Conjecture (1943)
If $G$ is loopless and has no $K_{t}$ minor then its chromatic number is $<t$

- Theorem The case $t=5$ is true (4CT proven 1976)
- Theorem The case $t=6$ is true (proven in 1993)
- $t>6$ is open, but: Theorem The conjecture is almost always true:


## Hadwiger's Conjecture is True for Almost Every Graph

B. Bollobás, P. A. Catlin* and P. Erdös

The contraction clique number $\operatorname{ccl}(G)$ of a graph $G$ is the maximal $r$ for which $G$ has a subcontraction to the complete graph $K^{r}$. We prove that for $d>2$, almost every graph of order $n$ satisfies $n\left(\left(\log _{2} n\right)^{\frac{1}{2}}+4\right)^{-1} \leqslant \operatorname{ccl}(G) \leqslant n\left(\log _{2} n-d \log _{2} \log _{2} n\right)^{-\frac{1}{2}}$. This inequality implies the statement in the title.

The bound need not to be sharp


- The Petersen graph contains a $K_{5}$ minor
- The Petersen graph is 3 -colorable

Thank you for your attention!

I hope that was of some help.

