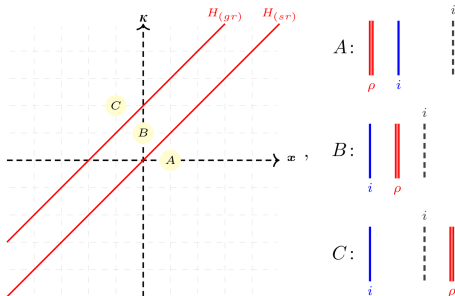


On weighted KLRW algebras

Or: Diagrammatic interpolation

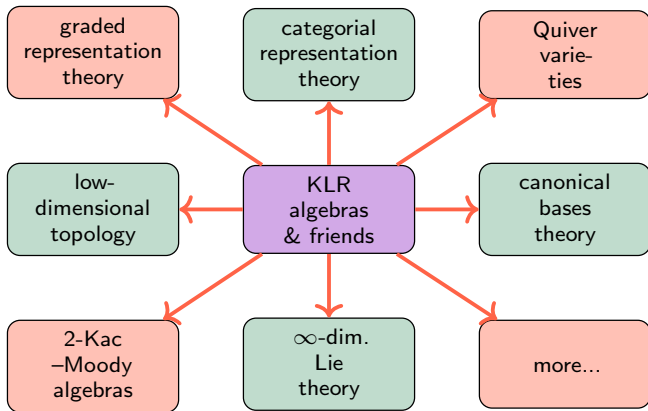
Daniel Tubbenhauer



Joint with Andrew Mathas

April 2022

Where are we?



- ▶ **Khovanov–Lauda–Rouquier ~2008 + many others (including BIT)**
KLR algebras are at the heart of categorical representation theory
- ▶ Similarly for quiver Schur algebras and diagrammatic Cherednik algebras
- ▶ **Problem** All of these are actually really complicated!

Observation

It often helps to find a “bigger” interpolating algebra, e.g.:

“Big”

$$\mathbb{C}[a, b][Y]/(Y - a)(Y - b)$$

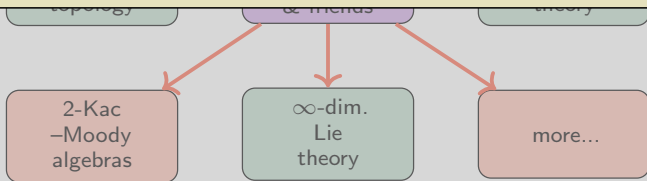
$a=0, b=0$

$a=1, b=-1$

“Small”

$$\mathbb{C}[Y]/(Y^2)$$

$$\mathbb{C}[Y]/(Y^2 - 1)$$



- ▶ **Khovanov–Lauda–Rouquier ~2008 + many others (including BIT)**
KLR algebras are at the heart of categorical representation theory
- ▶ Similarly for quiver Schur algebras and diagrammatic Cherednik algebras
- ▶ **Problem** All of these are actually really complicated!

Where

Observation

It often helps to find a “bigger” interpolating algebra, e.g.:

“Big”

$$\mathbb{C}[a, b][Y]/(Y - a)(Y - b)$$

$a=0, b=0$

$a=1, b=-1$

“Small”

$$\mathbb{C}[Y]/(Y^2)$$

$$\mathbb{C}[Y]/(Y^2 - 1)$$

Today

How to play the interpolation game using **planar geometry**?

As an upshot we get an algebra interpolating between various algebras appearing in categorical representation theory

The takeaway keyword: **Distance!**

- ▶ Kh
- ▶ KL
- ▶ Sim

(T)
ras

- ▶ **Problem** All of these are actually really complicated!

Where

Observation

It often helps to find a “bigger” interpolating algebra, e.g.:

“Big”

$$\mathbb{C}[a, b][Y]/(Y - a)(Y - b)$$

$a=0, b=0$

$a=1, b=-1$

“Small”

$$\mathbb{C}[Y]/(Y^2)$$

$$\mathbb{C}[Y]/(Y^2 - 1)$$

Careful

A lot but not everything

can be passed freely from $\mathbb{C}[a, b][Y]/(Y - a)(Y - b)$ to $\mathbb{C}[Y]/(Y^2)$
e.g. simple modules

The same is true for the algebras we see today!

Problem All of these are actually really complicated!

String diagrams – the baby case

Connect eight points at the bottom with eight points at the top:

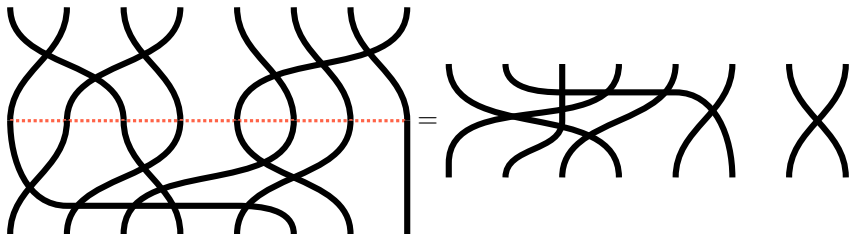


or



We just invented the symmetric group S_8 on $\{1, \dots, 8\}$

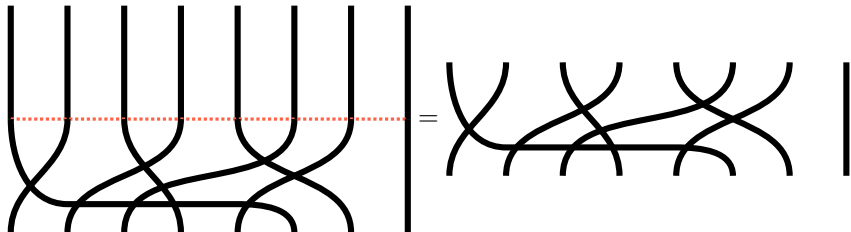
String diagrams – the baby case



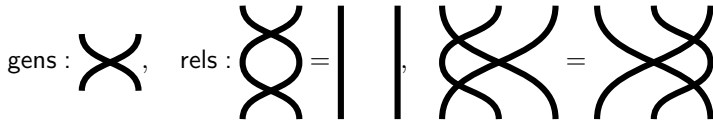
My multiplication rule for gh is “stack g on top of h ”

String diagrams – the baby case

- ▶ We clearly have $g(hf) = (gh)f$
- ▶ There is a do nothing operation $1g = g = g1$



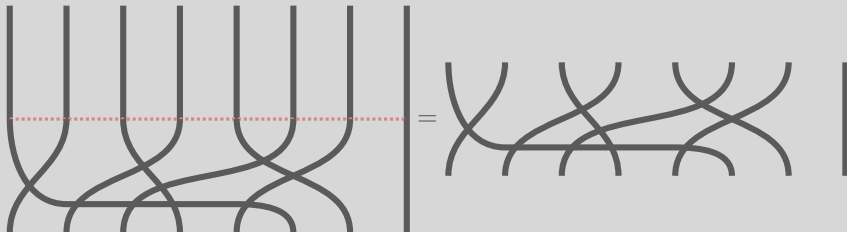
- ▶ Generators–relations (the Reidemeister moves)



The bait

In diagram algebras relations, properties, etc. become visually clear

- ▶ We clearly have
- ▶ There is a do nothing operation $1g = g = g1$



- ▶ Generators–relations (the Reidemeister moves)



String diagrams –

The bait

In diagram algebras relations, properties, etc. become visually clear

- ▶ We clearly have
- ▶ There is a do nothing operation $1g = g = g1$

The catch

Diagram algebras are usually “not really” using any planar geometry

For example, the diagrams for symmetric groups are just algebra written differently

- ▶ Generators–relations (the Reidemeister moves)



The bait

In diagram algebras relations, properties, etc. become visually clear

- ▶ We clearly have
- ▶ There is a do nothing operation $1g = g = g1$

The catch

Diagram algebras are usually “not really” using any planar geometry

For example, the diagrams for symmetric groups are just algebra written differently

- ▶ Generalization (the Braidification)
- ### Idea (Webster ~2012)

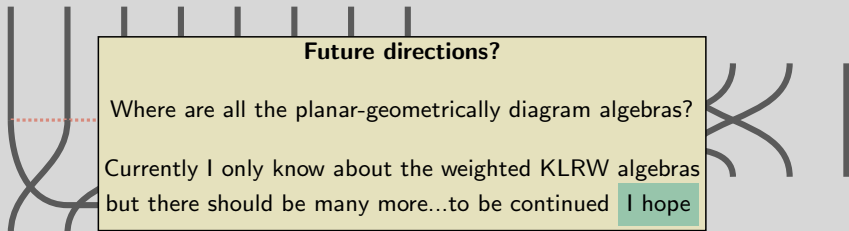
Define a diagram algebra that uses the distance in \mathbb{R}^2

The result is called **weighted KLRW algebra**

These are “planar-geometrically symmetric group diagram algebras”

String diagrams – the baby case

- ▶ We clearly have $g(hf) = (gh)f$
- ▶ There is a do nothing operation $1g = g = g1$



Future directions?

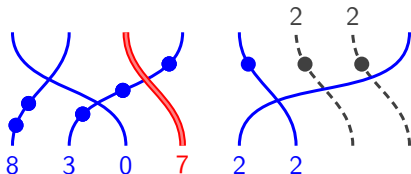
Where are all the planar-geometrically diagram algebras?

Currently I only know about the weighted KLRW algebras but there should be many more...to be continued **I hope**

- ▶ Generators–relations (the Reidemeister moves)



Weighted string diagrams



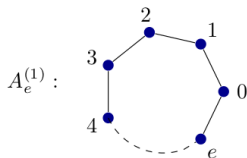
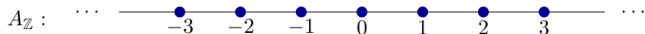
- ▶ Strings come in three types, **solid**, **ghost** and **red**

$$\text{solid} : \begin{array}{|c} \hline \\ \hline \end{array}, \quad \text{ghost} : \begin{array}{|c} \hline \cdots \\ \hline \end{array}, \quad \text{red} : \begin{array}{|c} \hline \\ \hline \end{array},$$

i *i* *i*

- ▶ Strings are labeled, and solid and ghost strings can carry dots
- ▶ Red strings **anchor** the diagram (red strings \leftrightarrow level)
- ▶ Otherwise no difference to symmetric group diagrams

Weighted string diagrams



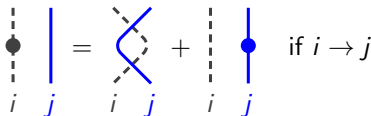
Examples of quivers Γ

An additional orientation fixes signs



- ▶ The strings are labeled by $i \in I$ from a fixed quiver $\Gamma = (I, E)$
- ▶ The relations (that I am not going to show you ;-)) depend on $e \in E$, e.g.:

“Reidemeister II
with error term” :



I usually never use the number π in a talk ;-)

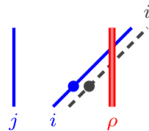
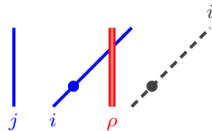
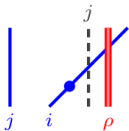
Weighted string diagrams

$$X = (-2\sqrt{3}, -\sqrt{2}, 0.5, \pi, 5) \longleftrightarrow \begin{array}{cccccc} | & | & || & | & | & | \\ -2\sqrt{3} & -\sqrt{2} & 0 & 0.5 & \pi & 5 \end{array}$$

Weighted quiver



diagram



- ▶ Choose endpoints $\mathbf{x} = (x_1, \dots, x_n) \in \mathbb{R}^n$, $\rho \in \mathbb{R}^\ell$ for the solid and red strings
- ▶ Choose a weighting $\sigma: E \rightarrow \mathbb{R}_{\neq 0}$ of the underlying graph $\Gamma = (I, E)$
- ▶ The weighted KLRW algebra **crucially depends** on these choices of endpoints! This is very different from “usual diagram algebras”

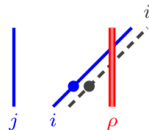
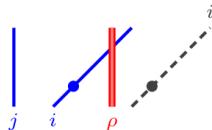
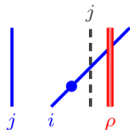
Weighted string diagrams

$$X = (-2\sqrt{3}, -\sqrt{2}, 0.5, \pi, 5) \iff \begin{array}{c} | \\ -2\sqrt{3} \end{array} \quad \begin{array}{c} | \\ -\sqrt{2} \end{array} \quad \begin{array}{c} || \\ 0 \end{array} \quad \begin{array}{c} | \\ 0.5 \end{array} \quad \begin{array}{c} | \\ \pi \end{array} \quad \begin{array}{c} | \\ 5 \end{array}$$

Weighted quiver



diagram



Weighting = ghost shifts

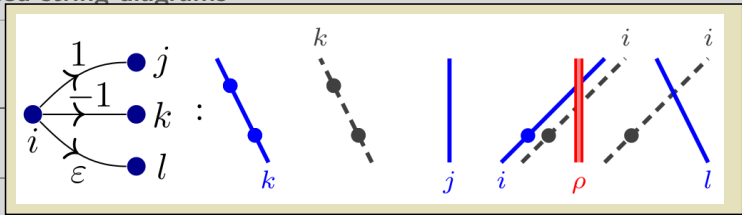
For $\epsilon: i \rightarrow j, \sigma_\epsilon > 0$, all solid i -strings get a ghost shifted $|\sigma_\epsilon|$ units and mimicking it
 For $\epsilon: i \rightarrow j, \sigma_\epsilon < 0$, all solid j -strings get a ghost shifted $|\sigma_\epsilon|$ units and mimicking it

► The weight of endpoints

This "asymmetric" definition, always shifting rightwards, makes life a bit more convenient but is not essential

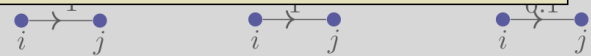
Weighted string diagrams

$X = (-$

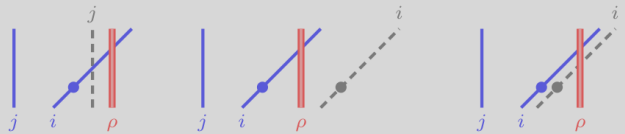


5

Weighted quiver



diagram



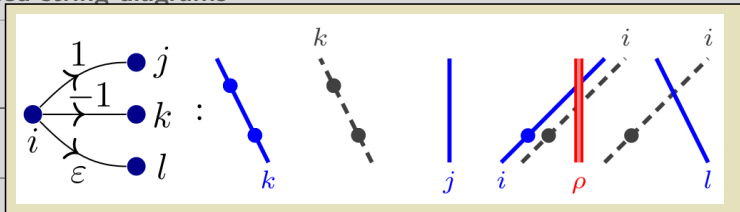
Weighting = ghost shifts

For $\epsilon: i \rightarrow j, \sigma_\epsilon > 0$, all solid i -strings get a ghost shifted $|\sigma_\epsilon|$ units and mimicking it
 For $\epsilon: i \rightarrow j, \sigma_\epsilon < 0$, all solid j -strings get a ghost shifted $|\sigma_\epsilon|$ units and mimicking it

- The weighted KLRW algebra **crucially depends** on these choices of endpoints! This is very different from "usual diagram algebras"

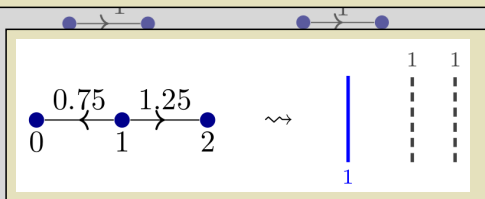
Weighted string diagrams

$X = (-$

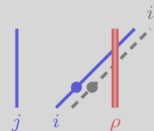


5

Weighted quiver



diagram



Weighting = ghost shifts

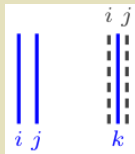
For $\epsilon: i \rightarrow j, \sigma_\epsilon > 0$, all solid i -strings get a ghost shifted $|\sigma_\epsilon|$ units and mimicking it
 For $\epsilon: i \rightarrow j, \sigma_\epsilon < 0$, all solid j -strings get a ghost shifted $|\sigma_\epsilon|$ units and mimicking it

- The weighted KLRW algebra **crucially depends** on these choices of endpoints! This is very different from “usual diagram algebras”

Weighted string diagrams

$$X = (-2\sqrt{3}, -\sqrt{2}, 0.5, \pi, 5) \leftrightarrow \begin{array}{cccccccc} | & & | & & || & | & & | & & | \\ \hline & & & & \circ & \circ & & & & \end{array}$$

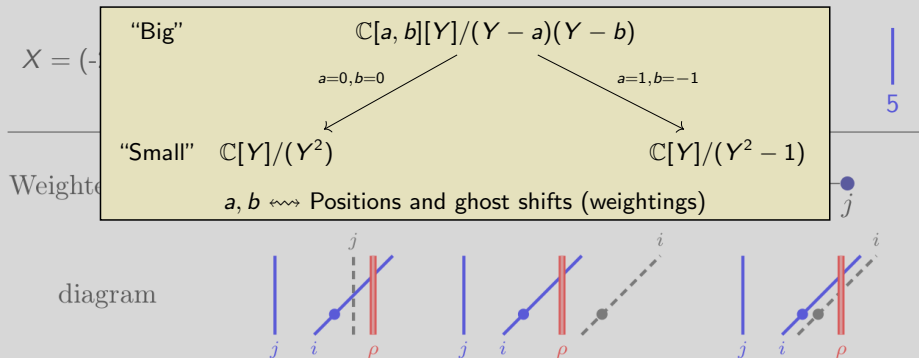
The following i and j -strings are not close:



Slogan Ghosts prevent the diagrams from being scale-able as for “usual diagram algebras”

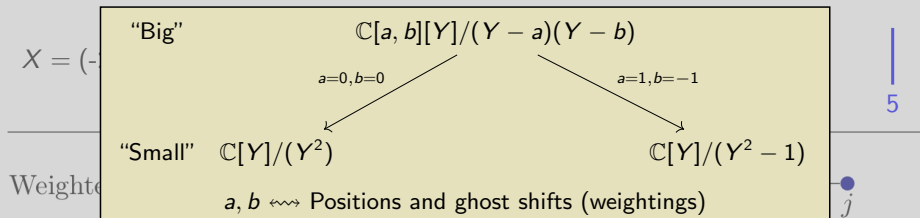
- ▶ Choose endpoints $\mathbf{x} = (x_1, \dots, x_n) \in \mathbb{R}^n$, $\rho \in \mathbb{R}^\ell$ for the solid and red strings
- ▶ Choose a weighting $\sigma: E \rightarrow \mathbb{R}_{\neq 0}$ of the underlying graph $\Gamma = (I, E)$
- ▶ The weighted KLRW algebra **crucially depends** on these choices of endpoints! This is very different from “usual diagram algebras”

Weighted string diagrams



- ▶ Choose endpoints $\mathbf{x} = (x_1, \dots, x_n) \in \mathbb{R}^n$, $\rho \in \mathbb{R}^\ell$ for the solid and red strings
- ▶ Choose a weighting $\sigma: E \rightarrow \mathbb{R}_{\neq 0}$ of the underlying graph $\Gamma = (I, E)$
- ▶ The weighted KLRW algebra **crucially depends** on these choices of endpoints! This is very different from “usual diagram algebras”

Weighted string diagrams



diagram

For “good choices” of X :

Semisimple	Huge ghost shifts
KLR	Tiny ghost shifts
Quiver Schur	Some specific “cluster” spacing
Diagrammatic Cherednik	Ghost shifts 1
Unnamed algebras	The rest

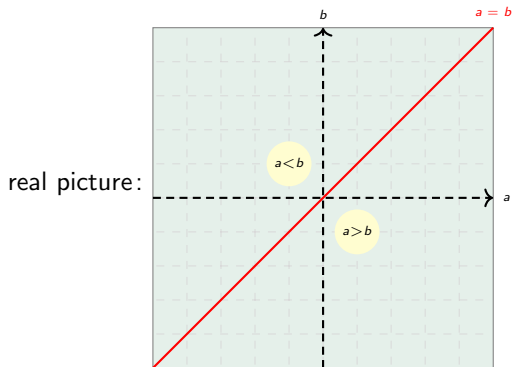
i

ρ

ed strings

- ▶ Choose
- ▶ Choose a weighting $\sigma: E \rightarrow \mathbb{R}_{\neq 0}$ of the underlying graph $\Gamma = (I, E)$
- ▶ The weighted KLRW algebra **crucially depends** on these choices of endpoints! This is very different from “usual diagram algebras”

Hyperplanes



- Specializations of $\mathbb{C}[a, b][Y]/(Y - a)(Y - b)$ come in two isomorphism classes:

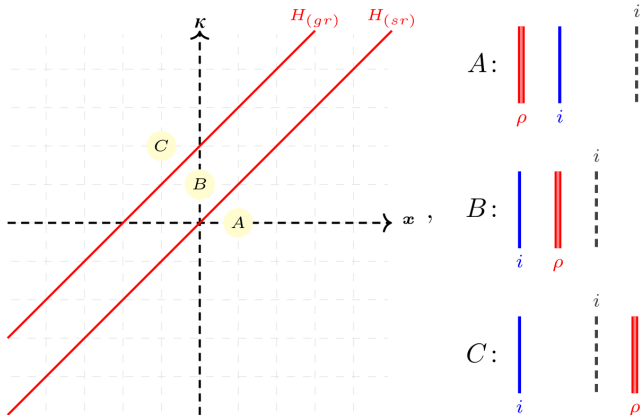
one double root $a = b$

&

two different roots $a \neq b$

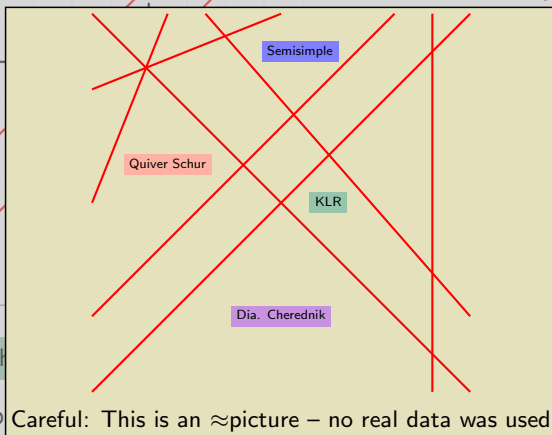
- What is the analog picture for weighted KLRW algebras?

Hyperplanes



- ▶ There is a **hyperplane arrangement (HA)** associated to weighted KLRW algebras
- ▶ The hyperplanes are defined by **“colliding strings”** (a form of distance)

- ▶ Alcoves of the HA \Rightarrow Morita equivalence classes of weighted KLRW algebras
 - ▶ There is a theory of translation functors
- ▶ \approx picture 1 There is an alcove for KLR, an alcove for the semisimple case etc.
 - ▶ \approx picture 2 Translation functors interpolate between these algebras



▶ There is a

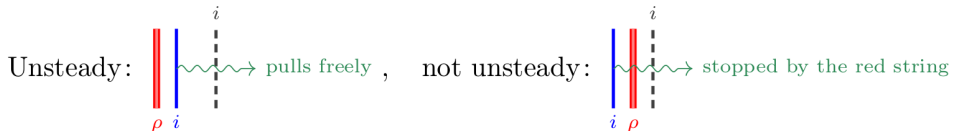
▶ The hyperp

KLRW algebras

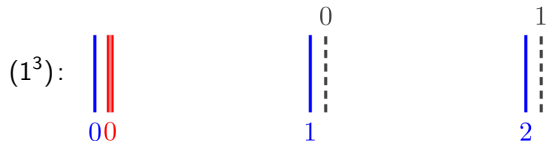
(f distance)

Distance is it!

- ▶ Cyclotomic (fin dim) quotients \Leftrightarrow bounded regions:



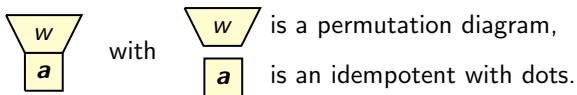
- ▶ Cellular bases \Leftrightarrow minimal regions (I will elaborate momentarily):



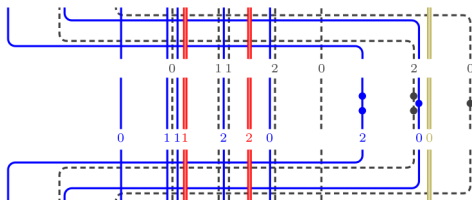
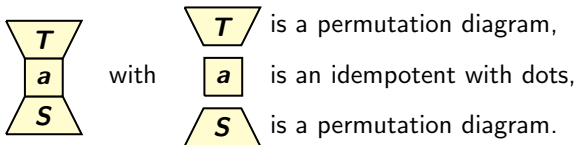
- ▶ More properties I won't explain today due to time restrictions...

Distance is it!

- ▶ Weighted KLRW algebras have **standard bases**, with the picture:



- ▶ Weighted KLRW algebras have **"cellular" bases**, with the picture:



Distance is

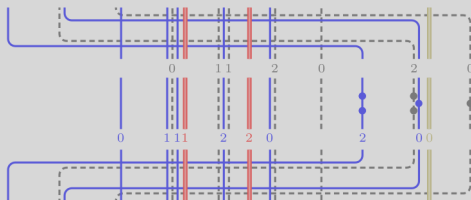
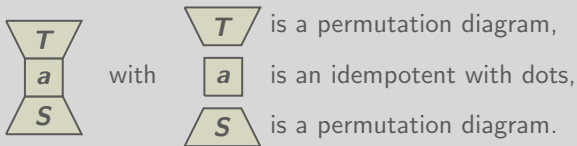
▶ Standard bases work regardless of the quiver but have no other property despite being a basis

▶ Weight

▶ Cellular bases depend on the quiver and give a classification of simple modules

(▶ Strictly speaking I should write “affine or sandwich” cellular but let us ignore that)

▶ Weighted KLRW algebras have “cellular” bases, with the picture:



Distance is

▶ Weight

▶ Standard bases work regardless of the quiver but have no other property despite being a basis

▶ Cellular bases depend on the quiver and give a classification of simple modules

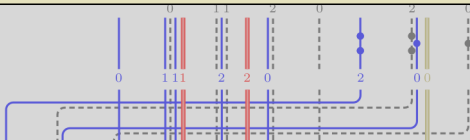
(▶ Strictly speaking I should write "affine or sandwich" cellular but let us ignore that)

▶ The overall strategy to construct cellular bases is the same for all types (but the details differ)

and for the infinite dimensional and the cyclotomic case the construction is also the same

▶ We know that the cellular bases work in types $A_{\mathbb{Z}}$, $A_e^{(1)}$, $B_{\mathbb{N}}$, $C_e^{(1)}$, $A_{2e}^{(2)}$, $D_{e+1}^{(2)}$ other, in particular finite, types are work in progress

▶ The combinatorics is inspired by, but different from, constructions of **Bowman** ~ 2017 , **Ariki–Park** $\sim 2012/2013$, **Ariki–Park–Speyer** ~ 2017



Distance is

▶ Weight

▶ Standard bases work regardless of the quiver but have no other property despite being a basis

▶ Cellular bases depend on the quiver and give a classification of simple modules

(▶ Strictly speaking I should write “affine or sandwich” cellular but let us ignore that)

▶ The overall strategy to construct cellular bases is the same for all types (but the details differ)

and for the infinite dimensional and the cyclotomic case the construction is also the same

▶ We know that the cellular bases work in types $A_{\mathbb{Z}}$, $A_e^{(1)}$, $B_{\mathbb{N}}$, $C_e^{(1)}$, $A_{2e}^{(2)}$, $D_{e+1}^{(2)}$ other, in particular finite, types are work in progress

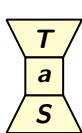
▶ The combinatorics is inspired by, but different from, constructions of **Bowman** ~ 2017 , **Ariki–Park** $\sim 2012/2013$, **Ariki–Park–Speyer** ~ 2017

I will now indicate how the construction works in type $C_e^{(1)}$

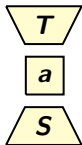
Why this type?

Because the code I am going to use works best for this type ;-)

Minimal diagrams in type $C_e^{(1)}$



with

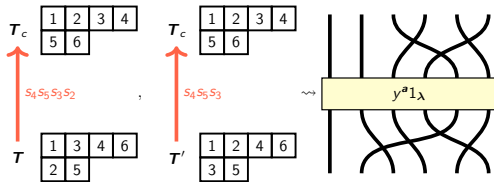


T is a permutation diagram,

a is an idempotent with dots,

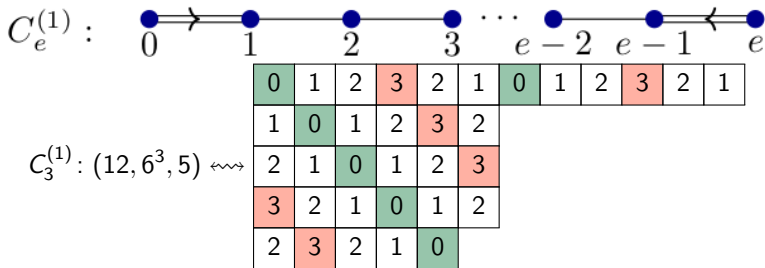
S is a permutation diagram.

- ▶ The definition of the permutation follows the usual strategy in this context:



- ▶ Let me focus on the middle $y^a 1_\lambda$

Minimal diagrams in type $C_e^{(1)}$



- ▶ Assume the tableaux combinatorics is given (a better statement later!)
- ▶ Place strings inductively **as far to the right as possible** (this is the order!)
- ▶ 1_λ is **minimal** with respect to placing the strings to the right
- ▶ 1_λ **stays minimal** when dots are put on certain strands \rightsquigarrow get $y^a 1_\lambda$
- ▶ **Done!**

Minimal diagrams in type $C_e^{(1)}$



Lets ignore the dots for today – I bothered you with too much combinatorics anyway ;-)
 But they come directly from the Reidemeister II relations, e.g.

for either of $\begin{cases} i = 0, j = 1, \\ i = e, j = e - 1, \end{cases}$

In other words: **Stare at Reidemeister II !**


- ▶ Place strings inductively **as far to the right as possible** (this is the order!)
- ▶ 1_λ is **minimal** with respect to placing the strings to the right
- ▶ 1_λ **stays minimal** when dots are put on certain strands \rightsquigarrow get $y^a 1_\lambda$
- ▶ **Done!**


Minimal diagrams in type $C_e^{(1)}$


$C_e^{(1)}$:


$C_3^{(1)}$: (12


Example for the middles $y^a 1_\lambda$

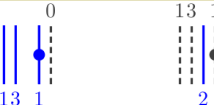
$y_{(1^5)} \mathbf{1}_{(1^5)} =$


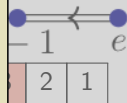
$y_{(3,1)} \mathbf{1}_{(3,1)} =$


$y_{(2,2)} \mathbf{1}_{(2,2)} =$


$y_{(5)} \mathbf{1}_{(5)} =$


$y_{(2,1^3)} \mathbf{1}_{(2,1^3)} =$


$y_{(4,1)} \mathbf{1}_{(4,1)} =$




► Assume the t

► Place strings

► 1_λ is minimal

► 1_λ stays minimal

► Done!

ment later!)

s is the order!)

ht

get $y^a 1_\lambda$

Minimal diagrams in type $C_e^{(1)}$



The (conjectural) picture for all types

$$[\Lambda_1] \xrightarrow{1} [-\Lambda_1 + \Lambda_2] \xrightarrow{2} [-\Lambda_2 + \Lambda_3] \xrightarrow{3} [\Lambda_2 - \Lambda_3] \xrightarrow{2} [\Lambda_1 - \Lambda_2] \xrightarrow{1} [-\Lambda_1]$$

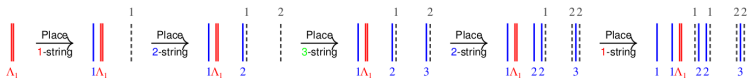


FIGURE 2. Top: A crystal in type B_3 . Bottom: Construction of an idempotent in a wKLRW algebra of type B_3 .

Checked for finite types (currently work in progress)

- ▶ 1_λ is minimal with respect to placing the strings to the right
- ▶ 1_λ stays minimal when dots are put on certain strands \rightsquigarrow get $y^a 1_\lambda$
- ▶ Done!

Minimal diagrams in type $C_e^{(1)}$



Wrap up

- ▶ Weighted KLRW algebras generalize KLR algebras and friends

- ▶ They have a built-in distance

- ▶ Most properties can be described using distance

- ▶ Most properties are type-independent

- ▶ Some properties are (in some form) type-independent

- ▶ Our dim calculations for the cellular basis match with the formulas of **Hu-Shi** ~2021 in the special cyclotomic KLR case

- ▶ 1_λ stays minimal when dots are put on certain strands \rightsquigarrow get $y^a 1_\lambda$

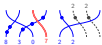
- ▶ Done!

Where are we?



- Khovanov-Lauda-Rouquier – 2006 + many others (including BFT)
- KLR algebras are at the heart of categorical representation theory
- Similarly for quiver Schur algebras and diagrammatic Chevalky algebras
- **Problem:** All of these are actually really complicated!

Weighted string diagrams



- Strings come in three types, **solid**, **ghost** and **red**.

solid : ghost : red :

- Strings are labeled, and solid and ghost strings can carry dots
- Red strings **cancel** the diagram (red strings \rightarrow level)
- Otherwise no difference to symmetric group diagrams

► Alluces of the BIA \rightarrow Morita equivalence classes of weighted KLRW algebras

- There is a theory of translation functors
- **Recipe 1:** There is an algebra for KLR, an algebra for the semisimple case etc.
- **Recipe 2:** Translation functors interpolate between these algebras.



- There is a **Caution:** This is an **recipe** – no real data was used
- The hypothesis is **Distance**

Where

Observation

It often helps to find a "bigger" interpolating algebra, e.g.:

"Big" $C[a, b][Y]/(Y^2 - aY - b)$

"Small" $C[Y]/(Y^2)$

How to play the interpolation game using **plectic geometry**?

As an upshot, we get an algebra interpolating between various algebras appearing in categorical representation theory.

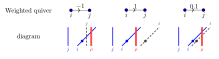
The takeaway keyword: **Distance**

- **KL**
- **KL**
- **KL**
- **Problem:** All of these are actually really complicated!

Weighted string diagrams

I usually never use the number c in a CBS :-)

$$X = (-2\sqrt{3}, -\sqrt{2}, 0.5, 0.5) \rightarrow \begin{matrix} 1 & 2 & 3 & 4 \\ -2\sqrt{3} & -\sqrt{2} & 0.5 & 0.5 \end{matrix}$$



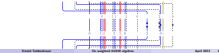
- Choose endpoints $x = (x_1, \dots, x_n) \in \mathbb{R}^n$, $p \in \mathbb{R}^n$ for the solid and red strings
- Choose a weighting $\alpha \in \mathbb{R} = \mathbb{R}_{\geq 0}$ of the underlying graph $F = (I, E)$
- The weighted KLRW algebra **usually depends** on those choices of endpoints! This is very different from "usual diagram algebra"

Distance is it!

- Weighted KLRW algebras have **standard bases**, with the picture:

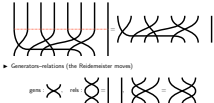


- Weighted KLRW algebras have **"cellular" bases**, with the picture:



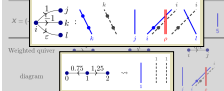
String diagrams – the baby case

- We clearly have $g(hf) = (gh)f$
- There is a **do nothing operation** $1g = g = g1$



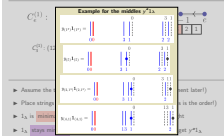
- **Generators-relations (the Reidemeister move)**

Weighted string diagrams



- **Weighting = ghost shifts**
- For $i \rightarrow j, i, j > 0$, all solid \rightarrow strings get a ghost shifted (\pm) units and mimicking in the $i \rightarrow j, i, j < 0$, all solid \rightarrow strings get a ghost shifted (\pm) units and mimicking in
- The weighted KLRW algebra **usually depends** on those choices of endpoints! This is very different from "usual diagram algebra"

Minimal diagrams in type $C_2^{(1)}$



- **Distance**

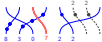
There is still much to do...

Where are we?



- Khovanov-Lauda-Rouquier – 2006 + many others (including BFT)
- KLR algebras are at the heart of categorical representation theory
- Similarly for quiver Schur algebras and diagrammatic Chevalky algebras
- **Problem:** All of these are actually really complicated!

Weighted string diagrams



- Strings come in three types, **solid**, **ghost** and **red**
- Solid strings are labeled, and solid and ghost strings can carry dots
- Red strings **anchor** the diagram (red strings \rightarrow level)
- Otherwise no difference to symmetric group diagrams

Algebras of the BIA \rightarrow Morita equivalence classes of weighted KLRW algebras

- There is a theory of translation functors
- **Proposition 1:** There is an algebra for KLR, an algebra for the semisimple case etc.
- **Proposition 2:** Translation functors interpolate between these algebras.



- There is a **Caution:** This is an **algebra** – no real data was used
- The hypothesis **Caution:** This is an **algebra** – no real data was used

Where

Observation

It often helps to find a "bigger" interpolating algebra, e.g.:

$$\begin{matrix} \text{"Big"} & C[\mathfrak{a}, \mathfrak{b}][Y]/(Y^2 - aY - b) \\ \swarrow \text{---} \text{---} \text{---} \searrow \\ \text{"Small"} & C[Y]/(Y^2) & C[Y]/(Y^2 - 1) \end{matrix}$$

How to play the interpolation game using **plectic geometry**?

As an upshot, we get an algebra interpolating between various algebras appearing in categorical representation theory

► **KL** \rightarrow $C[\mathfrak{a}, \mathfrak{b}][Y]/(Y^2 - aY - b)$

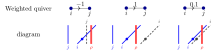
► **SL** \rightarrow $C[Y]/(Y^2 - 1)$

► **Problem:** All of these are actually really complicated!

Weighted string diagrams

I usually never use the number c in a CBS :-)

$$X = (-2\sqrt{3}, -\sqrt{2}, 0, 5) \rightarrow \begin{matrix} 1 & 1 & 1 & 1 \\ -2\sqrt{3} & -\sqrt{2} & 0 & 5 \end{matrix}$$



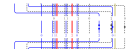
- Choose endpoints $x = (x_1, \dots, x_n) \in \mathbb{R}^n$, $p \in \mathbb{R}^n$ for the solid and red strings
- Choose a weighting $\mathfrak{a} = E = \mathbb{R}_{\geq 0}^m$ of the underlying graph $F = (I, E)$
- The weighted KLRW algebra is **generally diagrammatic** on these choices of endpoints! This is very different from "usual diagram algebra"

Distance is it!

- Weighted KLRW algebras have **standard bases**, with the picture:

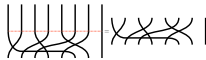


- Weighted KLRW algebras have **"cellular" bases**, with the picture:

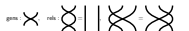


String diagrams – the baby case

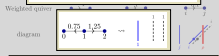
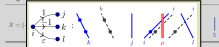
- We clearly have $g(hf) = (gh)f$
- There is a **do nothing operation** $1g = g = g1$



- Generators-relations (the Reidemeister moves)



Weighted string diagrams



- **Weighting = ghost shifts**
- For $i = j$, $a_i, a_j > 0$, all solid \rightarrow strings get a ghost shifted (a_i) units and mimicking in the $i = j$ $a_i, a_j < 0$, all solid \rightarrow strings get a ghost shifted (a_i) units and mimicking in
- The weighted KLRW algebra is **generally diagrammatic** on these choices of endpoints! This is very different from "usual diagram algebra"

Minimal diagrams in type $C_2^{(1)}$

Example for the minimal $y^2 \cdot 1_{\mathbb{Z}}$

$C_2^{(1)}$

$k_1 \otimes k_2 \otimes 1$	$\begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix}$	$\begin{pmatrix} 1 & 1 \\ 1 & 1 \end{pmatrix}$	$\begin{pmatrix} 1 & 1 \\ 1 & 1 \end{pmatrix}$
$k_1 \otimes k_2 \otimes 1$	$\begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix}$	$\begin{pmatrix} 1 & 1 \\ 1 & 1 \end{pmatrix}$	$\begin{pmatrix} 1 & 1 \\ 1 & 1 \end{pmatrix}$
$k_1 \otimes k_2 \otimes 1$	$\begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix}$	$\begin{pmatrix} 1 & 1 \\ 1 & 1 \end{pmatrix}$	$\begin{pmatrix} 1 & 1 \\ 1 & 1 \end{pmatrix}$

► Assume the $k_1 \otimes k_2 \otimes 1$ is the identity

► Place strings $\rightarrow 1_{\mathbb{Z}}$ in **matrix** (is the order)

► $1_{\mathbb{Z}}$ **matrix** (is the order)

► **Done!** get $y^2 \cdot 1_{\mathbb{Z}}$

Thanks for your attention!