

Lecture 5 Braided cats - Defs, examples, graphs

Def 5.1 A **braided** category \mathcal{C} is a \otimes -cat together with $\mathcal{T} = \{T_{X,Y} : X \otimes Y \xrightarrow{\sim} Y \otimes X\}$

such that

↳ a choice

Today
all \mathcal{C} 's
are braided

- naturality

$$T_{W,X}(fg) = (gf)T_{Y,Z}$$

- compatibility

$$T_{X,Y \otimes Z} = (\text{id}_Y \otimes T_{X,Z}) \circ (T_{X,Y} \otimes \text{id}_Z)$$

$$T_{X \otimes Y, Z} = (T_{X,Z} \otimes \text{id}_Y) \circ (\text{id}_X \otimes T_{Y,Z})$$

\mathcal{C} is further called

symmetric if
↳ a property

$$T_{Y,X} T_{X,Y} = \text{id}_{X \otimes Y}$$

$$\forall X, Y \in \mathcal{C}$$

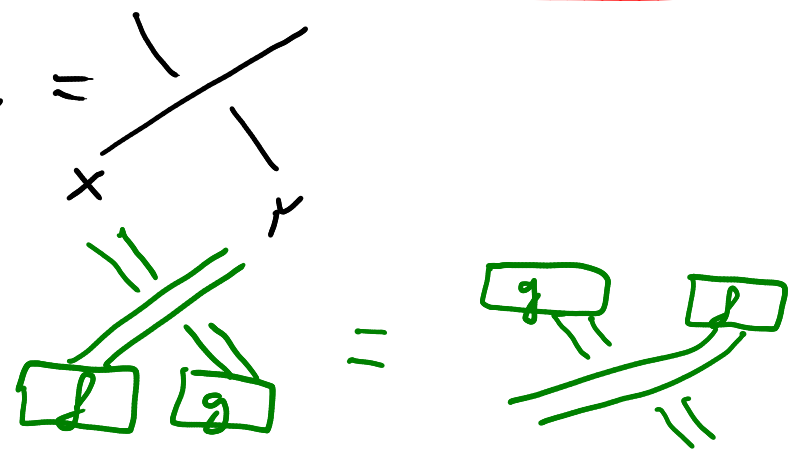
Def 5.2 A **braided** functor $F: \mathcal{C} \rightarrow \mathcal{D}$ is a functor such that

$$F(T_{X,Y}^{\mathcal{C}}) = T_{F(X), F(Y)}^{\mathcal{D}}$$

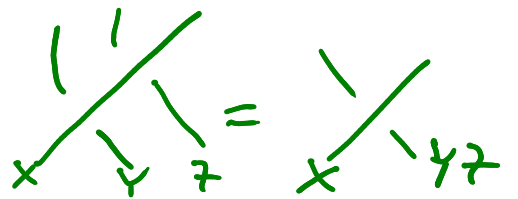
\leadsto We get the notion of equivalence of braided cats etc.

We draw:

\leadsto naturality



\leadsto compatibility



- sym. \leftarrow bad for knot theory :-)

Examples - (Sets, \times) is symmetric

- (Vect_K, \otimes) is symmetric

- $\text{Vect}^u(M)$ is braided iff M is commutative

(if M is not commutative, then $mn \neq nm$)

In case M is commutative, then braidings are in 1:1 correspondence with $c: M \times M \rightarrow K^*$ such that

$$c(g, h, j) = c(g, j) c(h, j) \omega(g, h, j)^{-1} \omega(h, g, j) \omega(h, j, g)^{-1}$$

$$c(g, h, j) = c(g, h) c(g, j) \omega(g, h, j) \omega(g, j, h)^{-1} \omega(j, g, h)$$

=> Classification of braidings (this is very hard beyond these toy cases)

Examples

- In particular, for $M = 1$, $\text{Vect}(1) \simeq \text{Vect}_{\mathbb{K}}$ allows only one braiding
- More general, if $w = 1$, then $c =$ bilinear pairing, which classifies the braidings
- These are symmetric iff c is symmetric
- If $M = \mathbb{Z}/2\mathbb{Z} = \langle 1, \sigma \rangle$, $w = 1$, then $c = 1$ defines a symmetric braiding.
- If $c(\sigma, \sigma) = -1$, then one gets a non-equivalent braiding and $(\text{Vect}(\mathbb{Z}/2\mathbb{Z}), \tau^c)$ is known as **super vector space**.

Examples

- The category Sym is the free symmetric cat. gen. by one object •

This is the braiding, of course

Recall: $Sym = \langle \cdot, \times \mid \chi = \parallel, \times = \times \rangle$

- The category TL only has an interesting braiding if linearized

- The category $Bv = \langle \cdot, \times, \cap, \cup \mid \chi = \parallel, \times = \times \rangle$ is symmetric and pivotal

$\cap = \cap, \times \cap = \times$
 $\cup = \cup + mirror$

Its quantum version qBv is braided and pivotal

Called twist

Bv is the free symmetric pivotal cat gen. by one self-dual object and trivial twist

From now on: \mathcal{C} is braided and pivotal

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Theorem 5.3 (Reidemeister, unoriented)

The free braided pivotal cat generated by one self dual object α has the presentation

$f\text{1Tan} = \langle \cdot, \times, \cup, \cap \mid \text{Reidemeister calc} \rangle$

$\rightarrow fR1: \text{two circles} = n = \text{one circle}$ "f-framed"

R2: $\text{cup} = \parallel$

R3: $\text{crossing} = \text{crossing}$

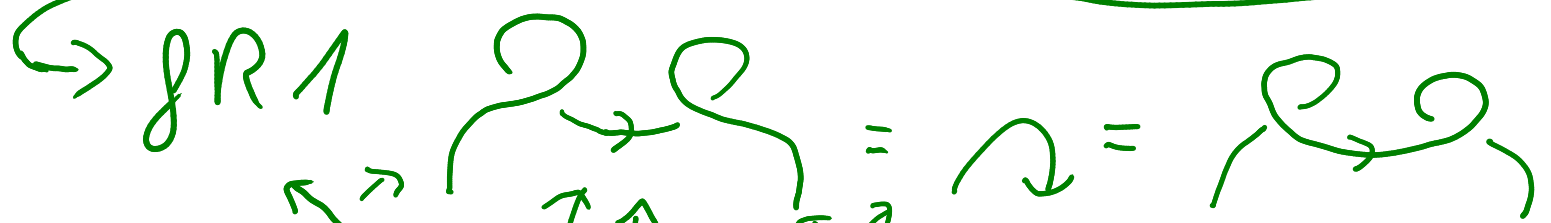
Iso: $\text{N} = \text{I} = \text{hook}$ $\text{hook} = \text{cap}$
 $\text{hook} = \text{cup}$

"framed 1-tangles"

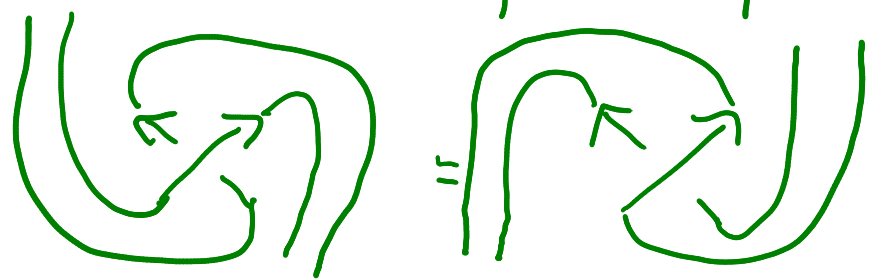
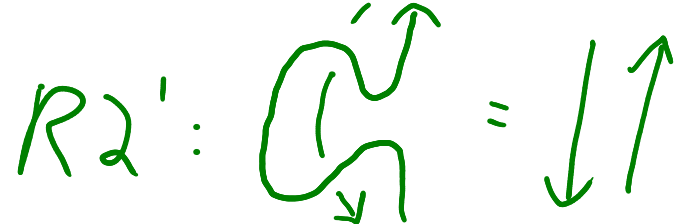
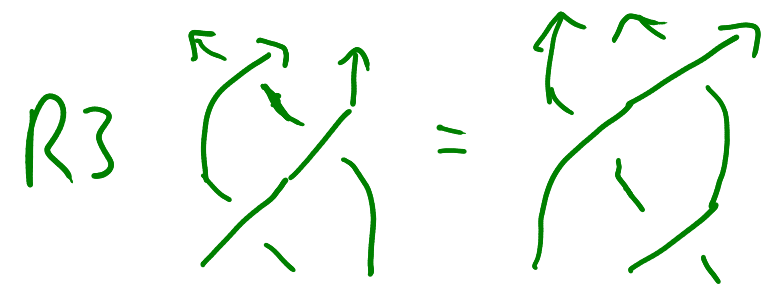
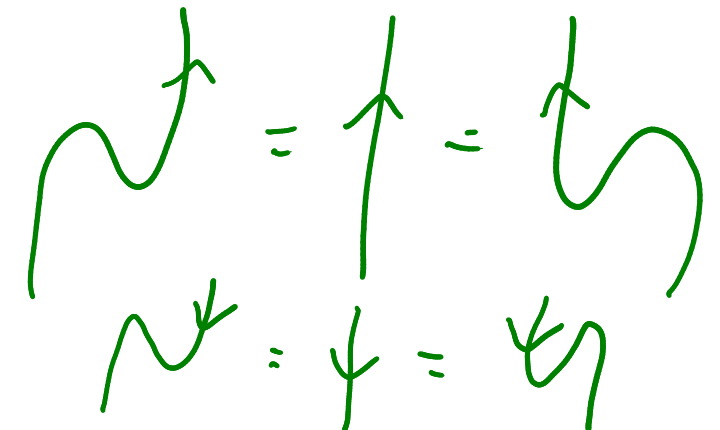
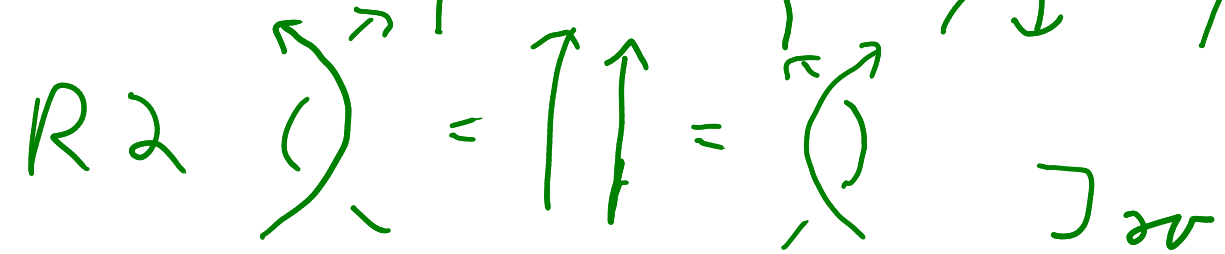
Theorem 5.4 (oriented version)

The free braided pivotal cat generated by one object + ($\rightsquigarrow \uparrow$) has the presentation

$\mathcal{F}1\text{Tan}^{\text{or}} = \langle +, \nearrow, \searrow, \cup, \cap, \cup, \cap \mid \text{or. Reidemeister} \rangle$



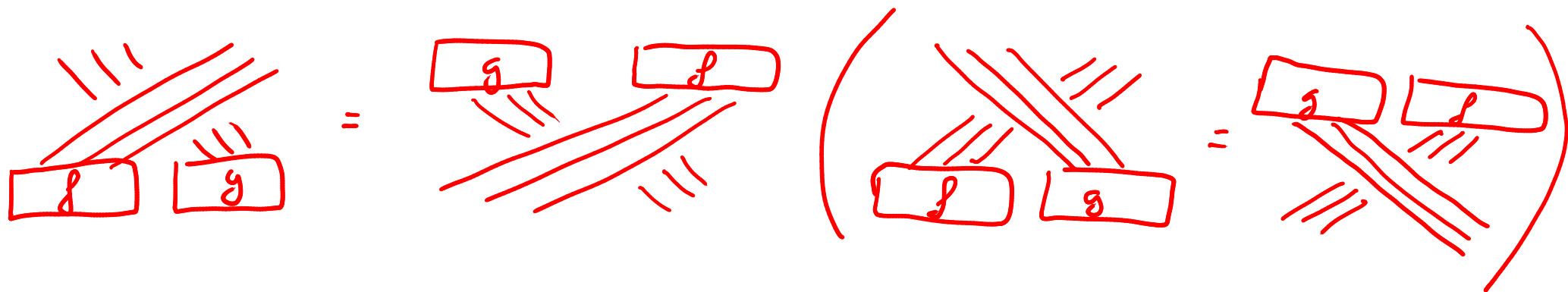
"oriented, framed 1-tangles"



Remark

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Formulated in other words, together with



Every braided pivotal category has a diagram calculus given by oriented, framed tangles

Of course, there are some consequences eg.

$$1 \circ 1 = \bigcirc = \sim = 1$$

Def 5.5

The *positive twists* Θ_x^+ are $\Theta_x^+ = \text{positive twist}$

The *negative twists* Θ_x^- are $\Theta_x^- = \text{negative twist}$

The halfmates of twists are $\tilde{\Theta}_x^+ = \text{halfmate of } \Theta_x^+$
 $\tilde{\Theta}_x^- = \text{halfmate of } \Theta_x^-$

Lemma 5.6

Twists and their halfmates are the same data

Proof.

$$\text{twist} = \text{halfmate}$$

$$\text{halfmate} = \text{twist}$$

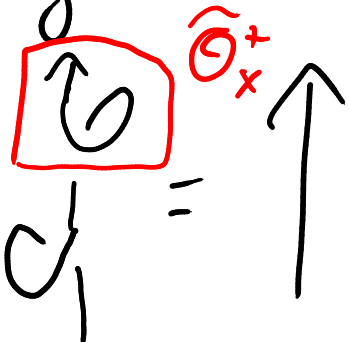
Lemma 5.7

$\tilde{\Theta}_x^+$ and $\tilde{\Theta}_x^-$ are natural isomorphism $id_x \Rightarrow id_x$

Moreover, $(\hat{\Theta}_x^+)^* = \tilde{\Theta}_{x^*}^-$ and $(\hat{\Theta}_x^-)^* = \tilde{\Theta}_{x^*}^+$

Proof

By the above, we have the equa. $\begin{matrix} \uparrow \\ \boxed{\circ} \\ \downarrow \\ \tilde{\Theta}_x^+ \end{matrix} = \uparrow$ and



$$\text{Now } \begin{matrix} \circ \\ \downarrow_x \end{matrix} = \cup \begin{matrix} \circ \\ \downarrow_x \end{matrix} = \downarrow_x = \begin{matrix} \circ^* \\ \downarrow \end{matrix} = \begin{matrix} \circ \\ \downarrow \end{matrix} = \begin{matrix} \circ^- \\ \downarrow_x \end{matrix}$$

Similarly for Θ_x^- .

Examples

- In $\text{fdVect}_{\mathbb{K}}$ the pairing is symmetric, thus $\overset{\uparrow}{\rho} = \underset{\downarrow}{\rho} = \overset{\uparrow}{\rho}^{\mathbb{M}}$

- Recall that for $G = \mathbb{Z}/3\mathbb{Z} = \langle 1, \sigma, \sigma^2 \rangle$ we can choose a pivotal structure

$$\overset{\sigma}{\curvearrowright} = 1 \quad \text{and} \quad \overset{\sigma}{\curvearrowleft} = \xi \quad \text{and} \quad \overset{\uparrow}{U}_{\sigma} = 1 \quad \overset{\downarrow}{U}_{\sigma} = \xi^2$$

Thus, for the trivial braiding we have

$$\overset{\uparrow}{\rho}_{\sigma} = \xi \underset{\downarrow}{\rho}_{\sigma} \quad \text{and} \quad \underset{\downarrow}{\rho}_{\sigma} = \overset{\uparrow}{\rho}_{\sigma} = \xi^2 \underset{\downarrow}{\rho}_{\sigma}$$

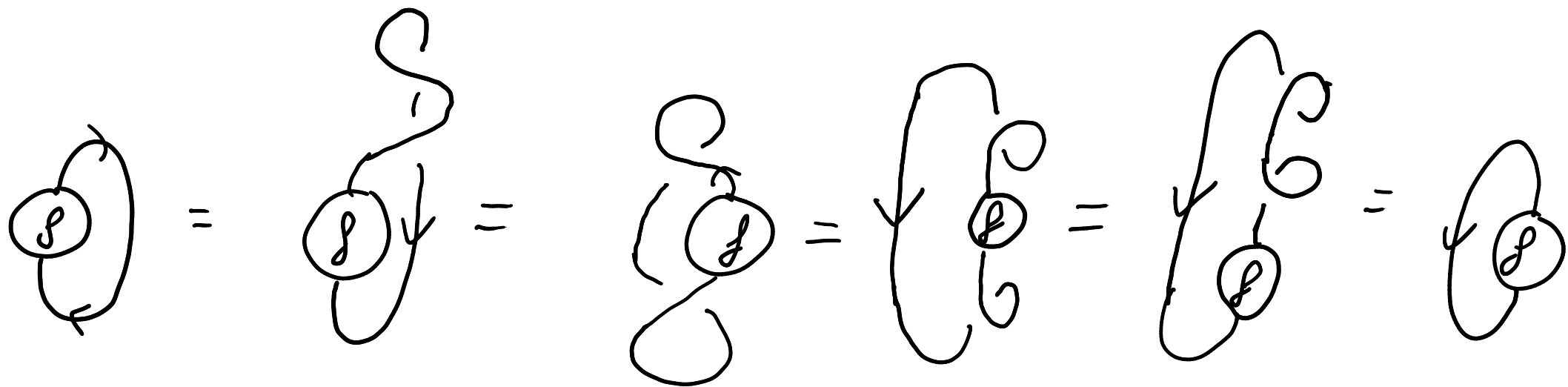
\Rightarrow twist are not trivial and also not equal
However $\overset{\uparrow}{\rho} = \underset{\downarrow}{\rho}$ is an important property motivating:

Def 5.8 \mathcal{C} is called **ribbon** if $\hat{\cup}_x = \cup_x \vee_x$ $1d$

lemma 5.9

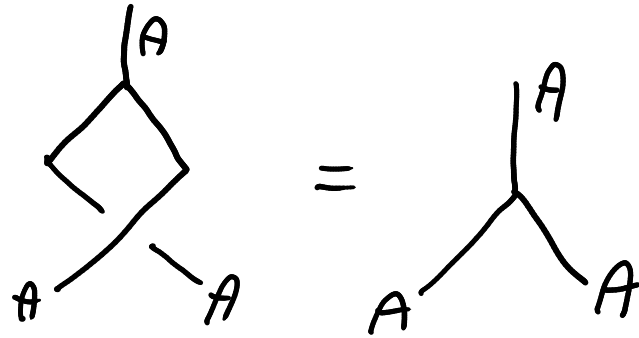
Ribbon categories are spherical

Proof



Here is the generalization of a comm. algebra 13

Def. 5.10 An algebra $A \in \mathcal{C}$ needs only braided is called **commutative** if



Similarly for coalgebras.

Examples

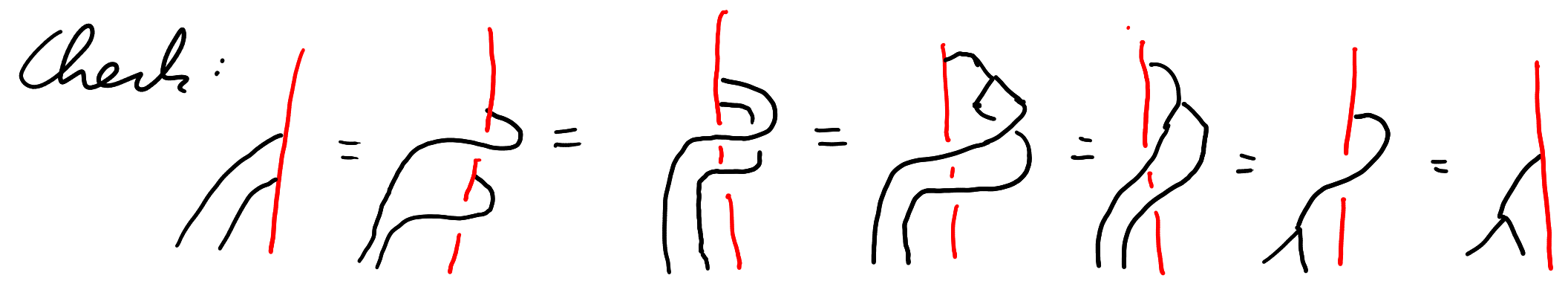
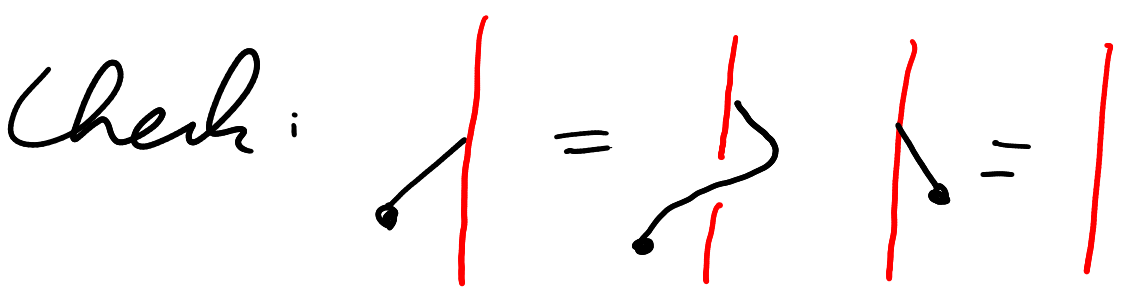
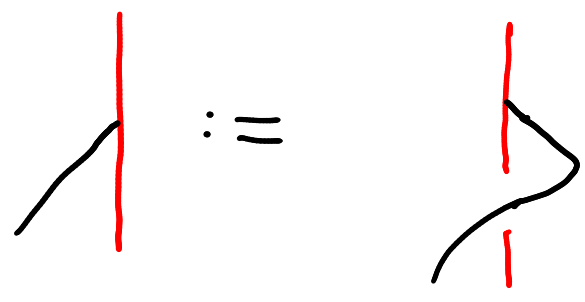
- In $\text{Vect}_{\mathbb{K}}$ comm algebras are comm algebras
- In $\text{Vect}(\mathbb{Z}/2\mathbb{Z})$ with the super braiding, comm algebras are supercommutative algebras, i.e. $\mathbb{Z}/2\mathbb{Z}$ -graded algebras such that $x \cdot y = (-1)^{\deg x \cdot \deg y} y \cdot x$

Proposition 5.11 Let $M \in \mathcal{C}$ be a right A -module object in $\mathcal{C} \leftarrow$ only braided $\hat{\mathcal{C}}$ comm. alg

Then M is also a right A module object.

In particular, $\text{Mod}_{\mathcal{C}}(A) \cong (A) \text{Mod}_{\mathcal{C}} \cong (A) \text{Mod}_{\hat{\mathcal{C}}}(A)$

Proof: Define



This implies that $\text{Mod}_{\mathcal{C}}(A)$ is a \otimes -cat.

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One can say more:

Proposition 5.12 The full subcategory $\text{Mod}_{\mathcal{C}}^{\circ}(A) \subset \text{Mod}_{\mathcal{C}}(A)$ given by A -modules satisfying

$$\begin{array}{c} | \\ \diagdown \\ | \\ \diagup \\ | \end{array} = \begin{array}{c} | \\ \diagdown \\ | \\ \diagup \\ | \end{array}$$

is braided with braiding inherited from \mathcal{C} .

Proof: Diagram considerations

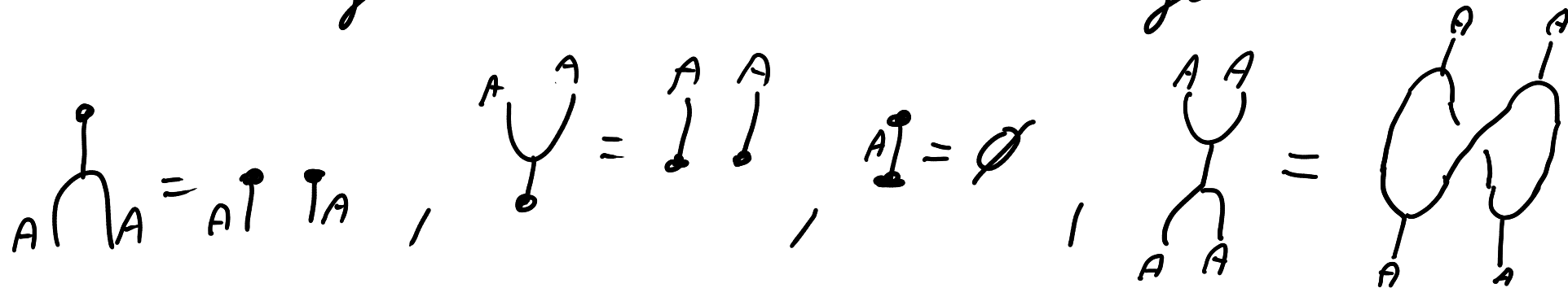
Example

If \mathcal{C} is symmetric, then $\text{Mod}_{\mathcal{C}}^{\circ}(A) = \text{Mod}_{\mathcal{C}}(A)$, which is thus also symmetric.

We can also generalize bialgebras

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Def 5.13 A **bialgebra** $A \in \mathcal{C} \leftarrow$ **only braided** is an algebra in \mathcal{C} and a coalgebra in \mathcal{C} such that

$$A \cap A = A \uparrow \downarrow A, \quad A \cup A = \downarrow \downarrow \uparrow \uparrow, \quad A \downarrow = \emptyset, \quad A \uparrow = \emptyset$$


Examples

- Bialgebras in $\text{Vect}_{\mathbb{K}}$ are bialgebras
- Bialgebras in $\text{Vect}(\mathbb{Z}/2\mathbb{Z})$ with the super braiding are super bialgebras

There is the evident notion of a bialgebra morphism etc.

Proposition 5.14 Let $B \in \mathcal{C}$ be a bialgebra

The category $\text{Mod}_e(B)$ is monoidal

Proof

Here is a way to define MN for $M, N \in \text{Mod}_e(B)$:

$$\begin{array}{c} | \\ \hline M \end{array} \otimes \begin{array}{c} | \\ \hline N \end{array} = \begin{array}{c} || \\ \hline MN \end{array}, \text{ action } \begin{array}{c} | \\ \hline \downarrow \end{array}$$

One can now check using the usual yoga, that this defines a \otimes -structure.