Formalize, Naturally!

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From the announcement of Formalize!(?):

... The derivation indicator view says that all proofs stand in some relation to a derivation, i.e. a mechanically checkable syntactical objects following fixed rules, that would not have any gaps.

. . .

Interactive and automated theorem provers promise to make the construction of a justification without any gaps feasible for complex mathematics.

. . .

Is this promise justified? Will the future of mathematical practice shift to more formal mathematics? Should it? ...

Jody Azzouni, 2004: The Derivation-Indicator View of Mathematical Practice

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Naproche:

- Natural proof assistant
- partial implementation of the derivation indicator view
- perfectly natural proof text is transformed into a formal derivation in some proof calculus
- natural language processing, logical transformations and automatic theorem proving

Martin Aigner, Günter M. Ziegler, 1998: Proofs from THE BOOK

Euclid's Proof. For any finite set $\{p_1, ..., p_r\}$ of primes, consider the number $n = p_1 \ p_2 \cdots p_r + 1$. This n has a prime divisor p. Put p is not one of the p_i : otherwise p would be a divisor of n and of the product $p_1 \ p_2 \cdots p_r$, and thus also of the difference $n - p_1 \ p_2 \cdots p_r = 1$, which is impossible. So a finite set $\{p_1, ..., p_r\}$ cannot be the collection of all prime numbers.

Aigner and Ziegler:

Euclid's Proof. For any finite set $\{p_1,...,p_r\}$ of primes,

consider the number $n = p_1 p_2 \cdots p_r + 1$.

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But p is not one of the p_i :

otherwise

p would be a divisor of n and of the product $p_1 \ p_2 \cdots p_r$,

and thus also of the difference $n-p_1\,p_2\cdots p_r=1$,

which is impossible.

So a finite set $\{p_1,...,p_r\}$ cannot be the collection of *all* prime numbers.

ForTheL text, accepted by Naproche:

Theorem 1. (Euclid) \mathbb{P} is infinite.

Proof. Assume that r is a natural number and p is a sequence of length r and $\{p_1, ..., p_r\}$ is a subset of \mathbb{P} .

Consider $n = p_1 \cdots p_r + 1$.

Take a prime divisor q of n.

Let us show that q is not equal to p_i for all i such that $1 \le i$ and $i \le r$.

Assume the contrary. Take i such that $1 \le i$ and $i \le r$ and $q = p_i$.

q is a divisor of n and q is a divisor of $p_1\cdots p_r$ (by 1).

Thus q divides 1.

Contradiction. qed.

Hence $\{p_1, ..., p_r\}$ is not the class of prime natural numbers. \Box

The Natural Proof Assistant Naproche

- (Original mathematical text)
- Input text in ForTheL (Formula Theory Language)
- Natural language and familiar symbolic terms
- LATEX format allows mathematical typesetting
- Natural language processing into first-order text
- Logical processing cuts up text into proof tasks
- Proof tasks are given to Automatic Theorem Prover (eprover)
- Eprover searches for superposition proofs
- (Eprover outputs derivations in superposition calculus)
- (Partial derivations can be combined into a complete derivation of the original text)

Naproche (Natural Proof Checking)

- Evidence Algorithm (Victor Glushkov, \sim 1970), ForTheL Input Language (Konstantin Vershinin, \sim 1975), System for Automated Deduction (Andrei Paskevich, \sim 2007)
- Naproche Project (Bernhard Schröder, PK, \sim 2002), (Old) Naproche System (Marcos Cramer, 2013)
- Naproche-SAD (Steffen Frerix, PK, 2018)
- Isabelle-Naproche (Steffen Frerix, Makarius Wenzel, PK, 2019): https://files.sketis.net/Isabelle_Naproche-20190611/

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Input

```
\begin{theorem} [Euclid]
$\Primes$ is infinite.
\end{theorem}
\begin{proof}
Assume that $r$ is a natural number and $p$ is a sequence of length $r$
and \left\{ \right\}  is a subset of \left\{ \right\} ?
Consider n=\Pr\{p\}\{1\}\{r\}+1\}.
Take a prime divisor $q$ of $n$.
Let us show that $q$ is not equal to $\val{p}{i}$ for all $i$ such that
1 \leq i \leq n and i \leq r.
Assume the contrary.
Take i such that 1 \leq i and i \leq r and q=\sqrt{p}{i}.
q is a divisor of n and q is a divisor of \Pr(q) = 1
1).
Thus $q$ divides $1$. Contradiction. qed.
Hence \int \{1\}\{r\} is not the class of prime natural numbers.
\end{proof}
```

Self-Contained Axiomatic Text (2-3 Pages LATEX Printout)

| [dump on] |
|---|
| Let $x \neq y$ stand for x is nonequal to |
| y. |
| [synonym number/-s] [synonym |

[synonym number/-s] [synonym divide/-s] [synonym set/-s] [synonym belong/-s] [synonym subset/-s]

1 Natural Numbers

Signature 2. A natural number is a notion.

Let i, k, l, m, n, p, q, r denote natural numbers.

Signature 3. 0 is a natural number.

Signature 4. 1 is a nonzero natural

. . .

number.

Signature 5. m+n is a natural number.

Signature 6. m*n is a natural number.

Axiom 7. m + n = n + m.

Axiom 8. (m+n) + l = m + (n+l).

2 The Natural Order

Definition 9. $m \le n$ iff there exists a natural number l such that m + l = n.

3 Division

4 Primes

Definition 10. n is prime iff n is nontrivial and for every divisor m of n m=1 or m=n.

Lemma 11. Every nontrivial k has a prime divisor.

Proof. Proof by induction.

5 Sets

Definition 12. \mathbb{N} is the class of natural numbers.

6 Sequences and Products

Signature 13. Let F be a sequence of length n such that $\{F_1, ..., F_n\} \subseteq \mathbb{N}$. $F_1 \cdots F_n$ is a nonzero natura number.

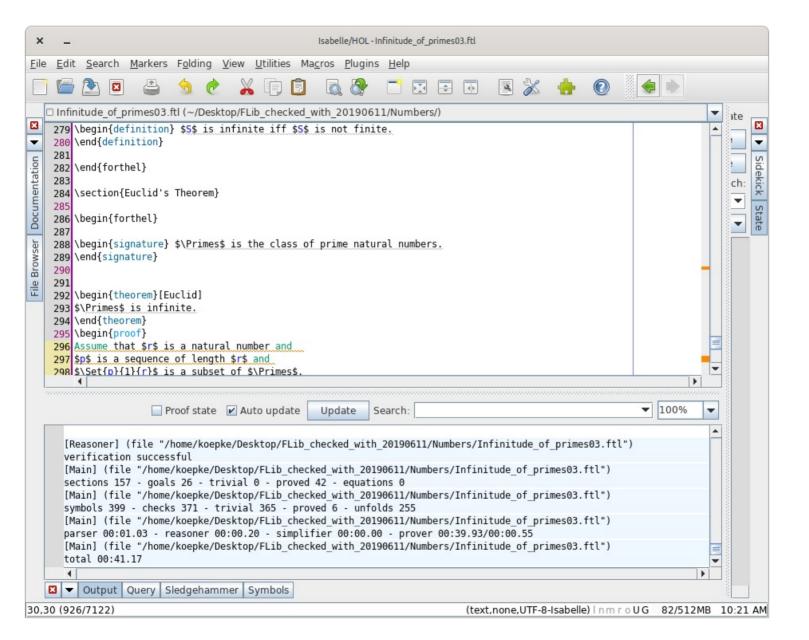
Finite and Infinite Sets

8 Euclid's Theorem

Signature 14. \mathbb{P} is the class of prime natural numbers.

Theorem 15. [Euclid] \mathbb{P} is infinite

Proof. ...



First-Order Translation

```
koepke@dell:~/TEST/Naproche-SAD$ stack exec Naproche-SAD -- -T ~/Desktop/
FLib_checked_with_20190611/Numbers/Infinitude_of_primes03.ftl
. . . . . .
hypothesis.
  assume forall v0 ((HeadTerm :: v0 = Primes) implies (aClass(v0) and forall v1
(aElementOf(v1,v0) iff (aNaturalNumber(v1) and isPrime(v1)))).
conjecture Euclid.
  isInfinite(Primes).
  proof.
    assume ((aNaturalNumber(r) and aSequenceOfLength(p,r)) and aSubsetOf(Set\{p\}\{1\}\{r\},
Primes)).
   n = Prod\{p\}\{1\}\{r\}+1.
    ((aNaturalNumber(q) and doDivides(q,n)) and isPrime(q)).
    forall v0 ((aNaturalNumber(v0) and (doLeq(1,v0) and doLeq(v0,r))) implies not q =
sdlbdtrb).
   proof.
      assume not thesis.
      (aNaturalNumber(i) and ((doLeq(1,i) and doLeq(i,r)) and q = sdlbdtrb)).
      ((aNaturalNumber(q) and doDivides(q,n)) and (aNaturalNumber(q) and doDivides(q,
Prod{p}{1}{r}))).
      doDivides(q,1).
      contradiction.
    ged.
    not (aClass(Set{p}{1}{r}) and forall v0 (aElementOf(v0,Set{p}{1}{r}) iff
(aNaturalNumber(v0) and isPrime(v0)))).
  qed.
```

Proof Goals, Sent to Eprover in TPTP Format

```
[Reasoner] (line 293 of "/home/koepke/Desktop/
FLib_checked_with_20190611/Numbers/Infinitude_of_primes03.ftl")
goal: Primes is infinite.
[Main] (file "/home/koepke/Desktop/FLib_checked_with_20190611/Numbers/
Infinitude_of_primes03.ftl")
fof(m_,hypothesis,$true).
fof(m_,hypothesis,aNaturalNumber(sz0)).
fof(m_,hypothesis,(aNaturalNumber(sz1) & ( ~ (sz1 = sz0)))).
fof(m_,hypothesis,( ! [W0] : ( ! [W1] : ((aNaturalNumber(W0) &
aNaturalNumber(W1)) => aNaturalNumber(sdtpldt(W0,W1))))).
fof(m_,hypothesis,( ! [W0] : ( ! [W1] : ((aNaturalNumber(W0) &
aNaturalNumber(W1)) => aNaturalNumber(sdtasdt(W0,W1))))).
fof(m_,hypothesis,( ! [W0] : ( ! [W1] : ((aNaturalNumber(W0) &
aNaturalNumber(W1)) => (sdtpldt(W0,W1) = sdtpldt(W1,W0))))).
```

Proof Goals (continued)

```
fof(m_,hypothesis,( ! [WO] : (aClass(WO) => (isInfinite(WO) <=> ( ~
isFinite(WO))))).
fof(m_,hypothesis,(aClass(szPzrzizmzezs) & ( ! [WO] : (aElementOf(WO,
szPzrzizmzezs) <=> (aNaturalNumber(W0) & isPrime(W0)))))).
fof(m__,conjecture,(( ! [WO] : ( ! [W1] : (((aNaturalNumber(W1) &
aSequenceOfLength(WO,W1)) & aSubsetOf(szSzeztlcdtrclcz1rclcdtrc(WO,
W1), szPzrzizmzezs)) => ( ? [W2] : ((W2 =
sdtpldt(szPzrzozdlcdtrclcz1rclcdtrc(W0,W1),sz1)) & ( ? [W3] :
(((aNaturalNumber(W3) & doDivides(W3,W2)) & isPrime(W3)) & ((!
[W4] : ((aNaturalNumber(W4) & (doLeq(sz1,W4) & doLeq(W4,W1))) =>
(~~(W3 = ssdlbdtrb(W0,W4)))) & (~~(![W4] : (aElementOf(W4,
szSzeztlcdtrclcz1rclcdtrc(W0,W1)) <=> (aNaturalNumber(W4) &
```

Proof Found!

```
[Main] (file "/home/koepke/Desktop/FLib_checked_with_20190611/Numbers/
Infinitude_of_primes03.ftl")
[eprover] # No SInE strategy applied
[Main] (file "/home/koepke/Desktop/FLib_checked_with_20190611/Numbers/
Infinitude_of_primes03.ftl")
[eprover] # Auto-Mode selected heuristic
G_E___208_C18_F1_SE_CS_SP_PS_S4Y
[Main] (file "/home/koepke/Desktop/FLib_checked_with_20190611/Numbers/
Infinitude_of_primes03.ftl")
[eprover] # and selection function SelectMaxLComplexAPPNTNp.
[Main] (file "/home/koepke/Desktop/FLib_checked_with_20190611/Numbers/
Infinitude_of_primes03.ftl")
[eprover] #
[Main] (file "/home/koepke/Desktop/FLib_checked_with_20190611/Numbers/
Infinitude_of_primes03.ftl")
[eprover] # Presaturation interreduction done
[Main] (file "/home/koepke/Desktop/FLib_checked_with_20190611/Numbers/
Infinitude_of_primes03.ftl")
[eprover] # Proof found!
[Main] (file "/home/koepke/Desktop/FLib_checked_with_20190611/Numbers/
Infinitude_of_primes03.ftl")
[eprover] # SZS status Theorem
```

Derivations generated by Eprover

```
[Main] (file "/home/koepke/Desktop/FLib_checked_with_20190611/Numbers/
Infinitude_of_primes03.ftl")
[eprover] cnf(c_0_47,hypothesis,(aElementOf(X1,
szPzrzizmzezs) | ~isPrime(X1) | ~aNaturalNumber(X1)),
inference(split_conjunct,[status(thm)],[c_0_39])).
[eprover] cnf(c_0_48, plain, (isPrime(esk29_2(X2,X1)))|^epred4_2(X1,X2)),
inference(split_conjunct,[status(thm)],[c_0_22])).
[eprover] cnf(c_0_49,plain,(aNaturalNumber(esk29_2(X2,
X1)) | ~epred4_2(X1,X2)), inference(split_conjunct,[status(thm)],
[c_0_22]).
[eprover] cnf(c_0_50,negated_conjecture,(~aElementOf(esk29_2(esk14_0,
esk15_0),szPzrzizmzezs)), inference(cn,[status(thm)],[inference(rw,
[status(thm)], [c_0_45, c_0_46])])).
[eprover] cnf(c_0_51,hypothesis,(aElementOf(esk29_2(X1,X2),
szPzrzizmzezs) | ~epred4_2(X2,X1)), inference(csr,[status(thm)],
[inference(spm, [status(thm)], [c_0_47, c_0_48]), c_0_49])).
[eprover] cnf(c_0_52, negated_conjecture, ($false), inference(cn,
[status(thm)], [inference(rw, [status(thm)], [inference(spm, [status(thm)],
[c_0_50, c_0_51]), c_0_46])]), ['proof']).
[eprover] # SZS output end CNFRefutation
```

Refutation Proof

```
R(\operatorname{sk}_{15}, \operatorname{sk}_{14})
X_{1} \in \mathbb{P} \vee \neg \operatorname{prime}(X_{1}) \vee \neg \operatorname{natural}(X_{1})
\operatorname{prime}(\operatorname{sk}_{29}(X_{2}, X_{1})) \vee \neg R(X_{1}, X_{2})
\operatorname{natural}(\operatorname{sk}_{29}(X_{2}, X_{1})) \vee \neg R(X_{1}, X_{2})
\operatorname{sk}_{29}(\operatorname{sk}_{14}, \operatorname{sk}_{15}) \notin \mathbb{P}
\operatorname{sk}_{29}(X_{1}, X_{2}) \in \mathbb{P} \vee \neg R(X_{2}, X_{1})
\bot
```

Statistics of Checking Euclid

- Text is checked in ≤ 1 minute
- 42 proofs found by eprover
- These proofs have between 2 and 106 clauses
- Project: combine such proofs to (superposition) derivations of complete texts
- Proof-checked ForTheL text ⇒ derivation
- Is that the derivation that the original proof in Aigner-Ziegler or Euclid "indicated"?
- _ ...

The Future of Formal Mathematics

- Interactive and automated theorem provers already allow the construction of justifications without any gaps for complex mathematics
- Mathematical practice will shift towards formal mathematics
- A decisive factor for the acceptance and momentum of this shift will be the ease and naturality of the interaction with the software
- Natural input languages are possible and will be provided for several interactive theorem provers
- Natural formal mathematics will require substantial further research and development

Thank You!