

Higher Set Theory - Classical and Ordinal Computability

Exercise Sheet 4
due on Tuesday, 3 May 2011

9. (a) Show that the remainder function $\text{Rem} : \omega \times \omega \rightarrow \omega$,

$$\text{Rem}(x, y) = \text{"the remainder of the division of } x \text{ by } y\text{"}$$

is primitive recursive.

- (b) Show that the Gödel β -function $\beta : \omega \times \omega \times \omega \rightarrow \omega$,

$$\beta(c, d, i) = \text{Rem}(c, 1 + (i + 1)d)$$

is primitive recursive.

- (c*) (*optional, no points awarded*) Show that for an arbitrary sequence $a_0, a_1, \dots, a_{n-1} \in \omega$ there are $c, d \in \omega$ such that for $i = 0, \dots, n - 1$ we have

$$\beta(c, d, i) = a_i.$$

Hint: Use the Chinese Remainder Theorem which states that for such $a_0, \dots, a_{n-1} \in \omega$ and pairwise relatively prime numbers $m_0, \dots, m_{n-1} \in \omega$ there is a number $c \in \omega$ s.t.

$$c \equiv a_i \pmod{m_i}$$

(4 points)

10. (a) Show that there is a register computable function $F : \omega \times \omega \rightarrow \omega$ s.t.

$$F(n, x) = F_n(x)$$

where $F_n : \omega \rightarrow \omega$ is the n -th unary primitive recursive function in some feasible enumeration.

- (b) Show, by a diagonal argument, that there is a register computable function that is not primitive recursive.

(8 points)