

Theorem 1. For every term G there is a term F satisfying:

$$G: V \rightarrow V \implies F: \text{Ord} \rightarrow V \wedge \forall \alpha F(\alpha) = G(F \upharpoonright \alpha).$$

Moreover for every term F' :

$$G: V \rightarrow V \wedge F': \text{Ord} \rightarrow V \wedge \forall \alpha F'(\alpha) = G(F' \upharpoonright \alpha) \implies F = F'.$$

Proof. Set

$$F = \{(\alpha, x) \mid \exists f: \alpha \rightarrow V \wedge \forall \xi < \alpha f(\xi) = G(f \upharpoonright \xi) \wedge x = G(f)\}.$$

Assume that $G: V \rightarrow V$.

(1) Let F', F'' be terms with

- $F': \text{dom}(F') \rightarrow V, F'': \text{dom}(F'') \rightarrow V,$
- $\text{dom}(F') \in \text{Ord}$ or $\text{dom}(F') = \text{Ord},$
- $\text{dom}(F'') \in \text{Ord}$ or $\text{dom}(F'') = \text{Ord},$
- $\forall \alpha \in \text{dom}(F') F'(\alpha) = G(F' \upharpoonright \alpha),$
- $\forall \alpha \in \text{dom}(F'') F''(\alpha) = G(F'' \upharpoonright \alpha).$

Then

$$\forall \alpha \in \text{dom}(F') \cap \text{dom}(F'') F'(\alpha) = F''(\alpha).$$

Proof. Assume not and let $\alpha \in \text{Ord}$ be minimal such that $\alpha \in \text{dom}(F') \cap \text{dom}(F'')$ and $F'(\alpha) \neq F''(\alpha)$. Then $\alpha \subseteq \text{dom}(F')$ and, by replacement, $F' \upharpoonright \alpha \in V$. Similarly, $\alpha \subseteq \text{dom}(F'')$ and $F'' \upharpoonright \alpha \in V$. By the minimality of α , $F' \upharpoonright \alpha = F'' \upharpoonright \alpha$. Hence

$$F'(\alpha) = G(F' \upharpoonright \alpha) = G(F'' \upharpoonright \alpha) = F''(\alpha).$$

Contradiction. *qed*(1)

Note that (1) proves the “moreover” part of the theorem.

(2) F is a function.

Proof. Let $(\alpha, x'), (\alpha, x'') \in F$ and take f', f'' satisfying $f': \alpha \rightarrow V \wedge \forall \xi < \alpha f'(\xi) = G(f' \upharpoonright \xi) \wedge x' = G(f')$ and $f'': \alpha \rightarrow V \wedge \forall \xi < \alpha f''(\xi) = G(f'' \upharpoonright \xi) \wedge x'' = G(f'')$. By (1), $f' = f''$ and hence

$$x' = G(f') = G(f'') = x''.$$

qed(2)

(3) $\text{dom}(F)$ is transitive.

Proof. Let $\alpha \in \text{dom}(F)$ and $\beta < \alpha$. Take f satisfying $f: \alpha \rightarrow V \wedge \forall \xi < \alpha f(\xi) = G(f \upharpoonright \xi) \wedge x = G(f)$. Then $f' = f \upharpoonright \beta$ satisfies $f': \beta \rightarrow V \wedge \forall \xi < \beta f'(\xi) = G(f' \upharpoonright \xi)$. $(\beta, G(f')) \in F$ and so $\beta \in \text{dom}(F)$. *qed*(3)

This implies immediately

(4) $\text{dom}(F) \in \text{Ord}$ or $\text{dom}(F) = \text{Ord}$.

(5) $\forall \alpha \in \text{dom}(F) F(\alpha) = G(F \upharpoonright \alpha)$.

Proof. Assume not and let $\alpha \in \text{Ord}$ be minimal such that $\alpha \in \text{dom}(F)$ and $F(\alpha) \neq G(F \upharpoonright \alpha)$. Set $f = F \upharpoonright \alpha$. By the minimality of α : $\forall \xi < \alpha f(\xi) = G(f \upharpoonright \xi)$. By the definition of F , $(\alpha, G(f)) \in F$ and

$$F(\alpha) = G(f) = G(F \upharpoonright \alpha),$$

contradiction. *qed*(5)

(6) $\text{dom}(F) = \text{Ord}$.

Proof. Assume not and let $\alpha = \text{dom}(F) \in \text{Ord}$. Set $f = F \upharpoonright \alpha$. By (5), $\forall \xi < \alpha f(\xi) = G(f \upharpoonright \xi)$. By the definition of F , $(\alpha, G(f)) \in F$ and $\alpha \in \text{dom}(F) = \alpha$, contradiction. $\square$