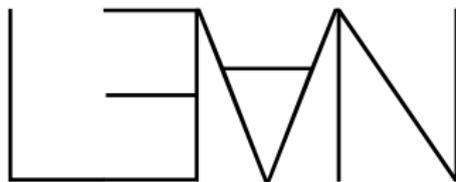


Tactics in Lean

Floris van Doorn

University of Pittsburgh

February 2, 2020



Mathlib contents

Subdirectory	LOC	Decls	Subdirectory	LOC	Decls
data	41849	10695	linear_algebra	4511	805
topology	17382	2709	computability	4205	575
tactic	12184	1679	group_theory	4191	1094
algebra	9830	2794	category	1770	389
analysis	7962	1237	number_theory	1394	228
order	6526	1542	logic	1195	403
category_theory	6299	1560	field_theory	1002	121
set_theory	6163	1394	geometry	848	70
measure_theory	6113	926	meta	784	135
ring_theory	5683	1080	algebraic_geometry	194	29

Total: 140k LOC (excluding whitespace/comments) and 34k declarations (as of December).

Kevin Buzzard (27 Sep 2017) I want to make these sorts of calculations trivial:

```
example : (((3 : ℝ)/4)-12)<6 := sorry
example : (6 : ℝ) + 9 = 15 := sorry
example : (2 : ℝ) * 2 + 3 = 7 := sorry
example : (5 : ℝ) ≠ 8 := sorry
example : (6 : ℝ) < 10 := sorry
example : (7 : ℝ)/2 > 3 := sorry
```

Mario Carneiro (2 Nov 2017) norm_num now solves all these goals.

example : $\neg (7-2)/(2*3) \geq (1:\mathbb{R}) + 2/(3^2)$:= by norm_num

example : $(6 : \mathbb{R}) + 9 = 15$:= by norm_num

example : $(2 : \mathbb{R})/4 + 4 = 3*3/2$:= by norm_num

example : $((3 : \mathbb{R})/4)-12 < 6$:= by norm_num

example : $(5 : \mathbb{R}) \neq 8$:= by norm_num

example : $(10 : \mathbb{R}) > 7$:= by norm_num

example : $(2 : \mathbb{R}) * 2 + 3 = 7$:= by norm_num

example : $(6 : \mathbb{R}) < 10$:= by norm_num

example : $(7 : \mathbb{R})/2 > 3$:= by norm_num

example : $(1103 : \mathbb{Z}) \leq (2102 : \mathbb{Z})$:= by norm_num

example : $(110474 : \mathbb{Z}) \leq (210485 : \mathbb{Z})$:= by norm_num

example : $(11047462383473829263 : \mathbb{Z}) \leq (21048574677772382462 : \mathbb{Z})$:=
by norm_num

example : $(210485742382937847263 : \mathbb{Z}) \leq (1104857462382937847262 : \mathbb{Z})$:=
by norm_num

example : $(210485987642382937847263 : \mathbb{N}) \leq (11048512347462382937847262 : \mathbb{N})$
:= by norm_num

example : $(210485987642382937847263 : \mathbb{Q}) \leq (11048512347462382937847262 : \mathbb{Q})$
:= by norm_num

[Kevin Buzzard](#) (11 Nov 2017) I love `norm_num`, I even use it to prove $0 < 1$ nowadays. I use it to prove everything. It's perfect. Many thanks for `norm_num`.

Johan Commelin (13 Mar 2019) It's quite humiliating, but how do I kill:

```
example (p : ℕ) [p.prime] : (p : ℝ) > 1 := sorry
```

Paul-Nicolas Madelaine (9 Apr 2019) Here is the first version of the cast tactic I've been working on.

```
example (a : ℕ) (b : ℤ) : (a : ℚ) < b ↔ (a : ℝ) < b := by norm_cast
example (a b : ℤ)      : a = b ↔ (a : ℚ) = b      := by norm_cast
example (a b : ℕ)      : (a : ℤ) + b = (a + b : ℕ) := by norm_cast
example (a : ℕ) (b : ℚ) : (a : ℂ) * b = ((a * b) : ℚ) := by norm_cast
example (a b : ℕ) : (((a : ℤ) : ℚ) : ℝ) + b = (a + (b : ℤ)) :=
by norm_cast
```

lift

Johan Commelin (9 Aug 2019) Suppose that $n : \mathbb{Z}$ and $h : n \geq 0$. Then every mathematician (and especially if they are new to Lean) wants to say $n : \mathbb{N}$. But that is not possible.

Floris van Doorn (10 Apr 2019) PR'd the `lift` tactic.

```
example {P : ℤ → Prop} (n : ℤ) (hn : n ≥ 0) : P n :=
begin
  lift n to ℕ using hn,
  /- New goal:
  P : ℤ → Prop,
  n : ℕ
  ⊢ P ↑n -/
  sorry
end
```

Other examples

Other useful tactics that have been implemented:

- `library_search` searches the library to close the current goal.
- `suggest` searches the library for a lemma that is applicable.
- `simp` `using` `h` closes the goal by simplifying both the goal and `h` to the same expression.
- `abel`, `ring`, `linarith`, `omega`: domain-specific automation.
- `tidy`, `finish`, `solve_by_elim`: general purpose automation .

rcases and rintro

`cases` destructs hypotheses, for example if $p : A \times B$ then `cases p with a b` gives two new hypothesis $a : A$ and $b : B$.

`rcases` and `rintro` perform these operations recursively.
Before:

```
cases h with y y2, cases y2 with yS hy, cases yS with y0 yx,
```

After:

```
rcases h with ⟨y, ⟨y0, yx⟩, hy⟩,
```

Before:

```
intro p, cases p with p1 p2, cases p1 with l hl, cases p2 with u hu,
```

After:

```
rintro ⟨⟨l, hl⟩, ⟨u, hu⟩⟩,
```

Before:

```
def yoneda C => (Cop => Type v1) :=
{ obj := λ X,
  { obj := λ Y, unop Y → X,
    map := λ Y Y' f g, f.unop >> g,
    /- (two fields omitted for readability) -/ },
  map := λ X X' f, { app := λ Y g, g >> f } }
```

@[simp] lemma obj_obj (X : C) (Y : C^{op}) :
 (yoneda.obj X).obj Y = (unop Y → X) := rfl

@[simp] lemma obj_map (X : C) {Y Y' : C^{op}} (f : Y → Y') :
 (yoneda.obj X).map f = λ g, f.unop >> g := rfl

@[simp] lemma map_app {X X' : C} (f : X → X') (Y : C^{op}) :
 (yoneda.map f).app Y = λ g, g >> f := rfl

After:

```
@[simps] def yoneda : C => (Cop => Type v1) :=
/- (definition is unchanged) -/
```

`#lint` is a semantic linter: it looks through the current file and looks for common mistakes in the declarations. Some mistakes that it catches:

- Have a hypothesis in a lemma that is never used;
- A declaration is incorrectly marked as a lemma or definition;
- A definition without documentation string.
- ...

localized notation

In Lean 3 notation is either local (to the current file or section) or global. You often want to use notation repeatedly, without it being global

```
localized "notation `ω` := ordinal.omega" in ordinal
```

You can get all notation in the ordinal locale by writing `open_locale ordinal`.

Writing Tactics: expr

We have reflection of expressions into Lean:

```
meta inductive expr (elaborated : bool := tt)
| var      {} : nat → expr
| sort    {} : level → expr
| const   {} : name → list level → expr
| mvar     : name → name → expr → expr
| local_const : name → name → binder_info → expr → expr
| app      : expr → expr → expr
| lam      : name → binder_info → expr → expr → expr
| pi       : name → binder_info → expr → expr → expr
| elet     : name → expr → expr → expr → expr
| macro    : macro_def → list expr → expr
```

For example,

```
λ (x : ℕ), nat.add x x
```

is reflected as

```
(lam x default (const nat []) (app (app (const nat.add []) (var 0)) (var 0)))
```

Writing Tactics: tactic

The tactic monad allows us to define custom tactics:

```
meta def tactic := interaction_monad tactic_state
```

A tactic ($t : \text{tactic } \alpha$) takes the current tactic state and runs a program to either

- succeed, and return the new tactic state and an element of α ;
- fail with an error message.

There are hooks for tactics implemented in C++:

```
meta constant infer_type : expr → tactic expr
```

assumption

This allows us to write our own tactics.

```
/-- `find_same_type t es` tries to find in `es` an expression with type  
    definitionally equal to `t` -/  
meta def find_same_type : expr → list expr → tactic expr  
| e []           := failed  
| e (H :: Hs) :=  
  do t ← infer_type H,  
    (unify e t >> return H) <|> find_same_type e Hs  
  
/-- `assumption` closes the goal if there is a hypothesis with the same type as  
    the goal. -/  
meta def assumption : tactic unit :=  
do { ctx ← local_context,  
    t   ← target,  
    H   ← find_same_type t ctx,  
    exact H }  
<|> fail "assumption tactic failed"  
  
example {p q : Prop} (h1 : p ∨ q) (h2 : q) : q := by assumption
```

Demo