

Frames for Mathematical Proofs

Workshop “Mathematical Language and Practical Type Theory”

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Outline

- 1 Introduction
- 2 Frames
- 3 Frames for Mathematical Texts
- 4 Further Frames
- 5 Frames and Mathematical Understanding
- 6 Conclusion

Motivation

Assumption: Different layers of interpretation of a mathematical text are useful at different stages of analysis and in different contexts.

Immediate Goal: make explicit in the formal representation of information that is implicit in the textual form

Theoretical Goal: bridge gap between formalist and textualist positions regarding mathematical proofs

Tools: from formal linguistics and artificial intelligence

Theses

- 1 FRAMES can serve as the basis for describing mathematical proofs.
- 2 Specifically, using FRAMES it is possible to model how mathematicians understand proofs that conform to proof patterns which have not been executed in a fully explicit way.
- 3 FRAMES can be used to model both (textual) structural properties of proofs and ontological aspects of mathematical knowledge. This distinction is useful.

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What are Frames?

Properties

- a concept in knowledge representation
↪ FILLMORE (1968) and MINSKY (1974)
- represent conceptual structure or prototypical situations
e.g. *birthday celebration, restaurant*.
- *roles* and *participants* (slots and fillers)
e.g. *waiter, diners, food, ...*
- organized in an *inheritance hierarchy*
typed feature structures (CARPENTER, 1992)

Usage

- e.g., in cognitive linguistics and artificial intelligence
 - explain how receiver completes information conveyed by sender
- ↪ linguistic project: FrameNet database (1,200 semantic frames)

Framing and Frames

One event can be framed differently, e.g. as *buying* and as *selling*

Frame: BUYING

[**BUYER** John] **bought** [**GOODS** a beautiful medieval book] [**TIME** yesterday].

Frame: SELLING

[**SELLER** Petra] **sold** [**GOODS** a beautiful medieval book] to [**BUYER** John] for [**MONEY** twenty Euros].

The Commercial Transaction frame from FrameNet

Frame Index

[A](#) [B](#) [C](#) [D](#) [E](#) [F](#) [G](#) [H](#) [I](#) [J](#) [K](#) [L](#) [M](#) [N](#) [O](#) [P](#) [Q](#) [R](#) [S](#)
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Commercial_transaction

[Lexical Unit Index](#)

Definition:

These are words that describe basic commercial transactions involving a **Buyer** and a **Seller** who exchange **Money** and **Goods**. The individual words vary in the frame element realization patterns. For example, the typical patterns for the verbs buy and sell are: BUYER buys GOODS from the SELLER for MONEY. SELLER sells GOODS to the BUYER for MONEY.

His **\$20** **TRANSACTION** **with Amazon.com** **for a new TV** had been very smooth.

FEs:

Core:

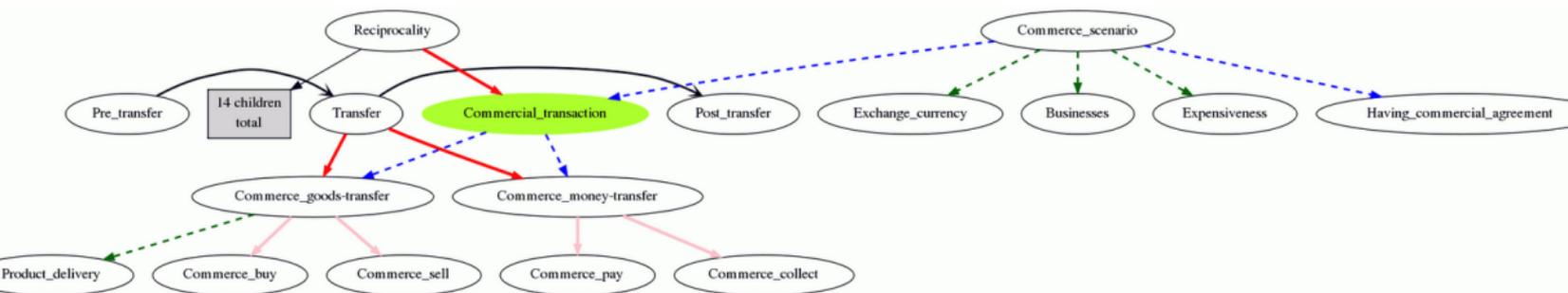
- Buyer [Bvr]** The **Buyer** wants the **Goods** and offers **Money** to a **Seller** in exchange for them.
- Goods [Gds]** The FE Goods is anything (including labor or time, for example) which is exchanged for Money in a transaction.
- Money [Mny]** Money is the thing given in exchange for Goods in a transaction.
- Seller [Slr]** The **Seller** has possession of the **Goods** and exchanges them for **Money** from a **Buyer**.

Non-Core:

- Means [Mns]** The means by which a commercial transaction occurs.
- Semantic Type:** State_of_affairs
- Rate [Rate]** Price or payment per unit of Goods.
- Unit [Unit]** The Unit of measure of the Goods according to which the exchange value of the Goods (or services) is set. Generally, it occurs in a by-PP.

Frame-frame Relations:

Frame-to-Frame relations: Multiple Inheritance



Screenshot <https://framenet.icsi.berkeley.edu/fndrupal/FrameGrapher>

Frames in Mathematical Texts

Goal: Model proofs and proof methods

Types of frames: (define types of slots)

Ontological: type of mathematical object

e.g. **Circle**, **slots:** center, radius, diameter, circumference, ...

e.g. **Vector Space**, **slots:** zero, unit, field, dimension, ...

Structural: part of proofs

e.g. **Induction**, **slots:** induction variable, hypothesis, step, domain,

...

e.g. **Extremal Proof**, **slots:** object type, initial object, parameter

Frames Example: Induction

Induction Frame (structural)

with slots:

- BASE CASE
- INDUCTION HYPOTHESIS
- INDUCTION STEP
- INDUCTION DOMAIN: **Inductive Type (ontological)** with
 - BASE CONSTRUCTOR
 - RECURSIVE CONSTRUCTOR
- Induction variable

(see FISSENI, SARIKAYA, SCHMITT and SCHRÖDER, 2019)

induction

INDUCTION-DOMAIN \boxed{d} $\left[\begin{array}{l} \textit{inductive-type} \\ \text{BASE-CONSTRUCTOR} \quad \boxed{bc} \textit{ base-constructor} \\ \text{RECURSIVE-CONSTRUCTOR} \quad \boxed{rc} \textit{ recursive-constructor} \end{array} \right]$

INDUCTION-VARIABLE $\left[\begin{array}{l} \textit{variable} \\ \text{NAME} \quad \boxed{x} \textit{ symbolic} \\ \text{TYPE} \quad \boxed{d} \end{array} \right]$

ASSERTION $\forall \boxed{x}. \boxed{S}$

induction-proof

INDUCTION-SIGNATURE $\left[\begin{array}{l} \textit{induction-signature} \\ \text{INDUCTION-HYPOTHESIS} \quad \boxed{ih} \textit{ sentence} \\ \text{STEP-FUNCTION} \quad \boxed{?!} \boxed{rc} \\ \text{BASE-CONDITION} \quad \boxed{bcc} \boxed{?!} \boxed{x} = \boxed{bc} \\ \text{INDUCTION-CONDITION} \quad \boxed{icc} \boxed{?!} \boxed{x} = \boxed{rc}(\dots) \end{array} \right]$

PROOF

BASE-CASE $\left[\begin{array}{l} \textit{proved-under-hypothesis} \\ \text{HYPOTHESIS} \quad \boxed{bcc} \\ \text{THESIS} \quad \boxed{S} \\ \text{ASSERTION} \quad \boxed{bcc} \Rightarrow \boxed{S} \\ \text{PROOF} \quad \textit{list}(\textit{proof-step} \vee \textit{assumption} \vee \textit{definition} \vee \textit{goal}) \end{array} \right]$

[...]

ASSERTION $\forall x. S$ *induction-proof*

INDUCTION-SIGNATURE

*induction-signature*INDUCTION-HYPOTHESIS $[ih]$ sentenceSTEP-FUNCTION $(?!)$ $[rc]$ BASE-CONDITION $[bcc]$ $(?!)$ $x = [bc]$ INDUCTION-CONDITION $[icc]$ $(?!)$ $x = [rc](...)$ PROOF

BASE-CASE

*proved-under-hypothesis*HYPOTHESIS $[bcc]$ THESIS $[S]$ ASSERTION $[bcc] \Rightarrow [S]$ PROOF *list(proof-step v assumption v definition v goal)*

INDUCTION-STEP

*proved-under-hypothesis*HYPOTHESIS $[icond]: ([icc] \wedge [ih])$ THESIS $[S]$ ASSERTION $[icond] \Rightarrow [S]$ PROOF *list(proof-step v assumption v definition v goal)*

An Induction Proof

KOWALSKI (2016, p. 93)

Proof. First, the second statement is indeed more precise than the first: let $k \geq 1$ be such that $f^k = 0$ but $f^{k-1} \neq 0$; there exists $v \neq 0$ such that $f^{k-1}(v) \neq 0$, and we obtain $k \leq n$ by applying the second result to this vector v . We now prove the second claim. Assume therefore that $v \neq 0$ and that $f^k(v) = 0$ but $f^{k-1}(v) \neq 0$. Let t_0, \dots, t_{k-1} be elements of K such that

$$t_1 v + \dots + t_{k-1} f^{k-1} f^{[sic]}(v) = 0.$$

Apply f^{k-1} to this relation; since $f^k(v) = \dots = f^{2k-2} f^{[sic]}(v) = 0$, we get

$$t_1 f^{k-1}(v) = t_1 f^{k-1}(v) + t_2 f^k(v) + \dots + t_{k-1} f^{2k-2} f^{[sic]}(v) = 0,$$

and therefore $t_1 f^{k-1}(v) = 0$. Since $f^{k-1}(v)$ was assumed to be non-zero, it follows that $t_1 = 0$.

Now repeating this argument, but applying f^{k-2} to the linear relation (and using the fact that $t_1 = 0$), we get $t_2 = 0$.

Then similarly we derive **by induction** that $t_i = 0$ for all i , proving the linear independence stated.

in the first equation, the exponent $k - 1$ has to be replaced by $k - 2$; in the line below and the second equation, $2k - 2$ by $2k - 3$.

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$$t_1 v + \dots + t_{k-1} f^{k-1}([sic!])(v) = 0.$$

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induction

INDUCTION-DOMAIN

inductive-type
 [BASE-CONSTRUCTOR $\boxed{\text{bc}}$ *base-constructor*
 RECURSIVE-CONSTRUCTOR $\boxed{\text{rc}}$ *recursive-constructor*]

INDUCTION-VARIABLE

variable
 [NAME $\boxed{\text{v}}$ *symbolic*
 TYPE $\boxed{\text{t}}$]

ASSERTION

$\forall \boxed{\text{v}} . \boxed{\text{a}}$

induction-proof

INDUCTION-SIGNATURE

induction-signature
 INDUCTION-HYPOTHESIS $\boxed{\text{ih}}$ *sentence*
 STEP-FUNCTION $\boxed{\text{sf}}$ (?)
 BASE-CONDITION $\boxed{\text{bcc}}$ (?) $\boxed{\text{a}}$ = $\boxed{\text{bc}}$
 INDUCTION-CONDITION $\boxed{\text{icc}}$ (?) $\boxed{\text{a}}$ = $\boxed{\text{rc}}(\dots)$

PROOF

BASE-CASE

proved-under-hypothesis
 HYPOTHESIS $\boxed{\text{bcc}}$
 THESIS $\boxed{\text{a}}$
 ASSERTION $\boxed{\text{bcc}} \Rightarrow \boxed{\text{a}}$
 PROOF *list(proof-step \vee assumption \vee definition \vee goal)*

INDUCTION-STEP

proved-under-hypothesis
 HYPOTHESIS $\boxed{\text{icand}}$: $(\boxed{\text{bcc}} \wedge \boxed{\text{ih}})$
 THESIS $\boxed{\text{a}}$
 ASSERTION $\boxed{\text{icand}} \Rightarrow \boxed{\text{a}}$
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PROOF

BASE-CASE

proved-under-hypothesis

HYPOTHESIS $\boxed{\&}$
 THESIS $\boxed{?}$
 ASSERTION $\boxed{\&} \Rightarrow \boxed{?}$
 PROOF $\text{list}(\text{proof-step} \vee \text{assumption} \vee \text{definition} \vee \text{goal})$

INDUCTION-STEP

proved-under-hypothesis

HYPOTHESIS $\boxed{\& \text{cond}} : (\boxed{\&} \wedge \boxed{\&})$
 THESIS $\boxed{?}$
 ASSERTION $\boxed{\& \text{cond}} \Rightarrow \boxed{?}$
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[<u>Q</u>]	BASE-CONSTRUCTOR	[<u>bc</u>]	base-constructor
	RECURSIVE-CONSTRUCTOR	[<u>rc</u>]	recursive-constructor

INDUCTION-VARIABLE

variable

[<u>s</u>]	NAME	symbolic
[<u>T</u>]	TYPE	

ASSERTION

[V]. [I]*induction-proof*

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[<u>ic</u>]	INDUCTION-CONDITION	(?) [<u>s</u>] = [<u>rc</u>](...)

PROOF

BASE-CASE

proved-under-hypothesis

[<u>h</u>]	HYPOTHESIS	[<u>bc</u>]
[<u>T</u>]	THESIS	[<u>s</u>]
[<u>A</u>]	ASSERTION	[<u>bc</u>] \Rightarrow [<u>s</u>]
[<u>P</u>]	PROOF	list(proof-step \vee assumption \vee definition \vee goal)

INDUCTION-STEP

proved-under-hypothesis

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PROOF

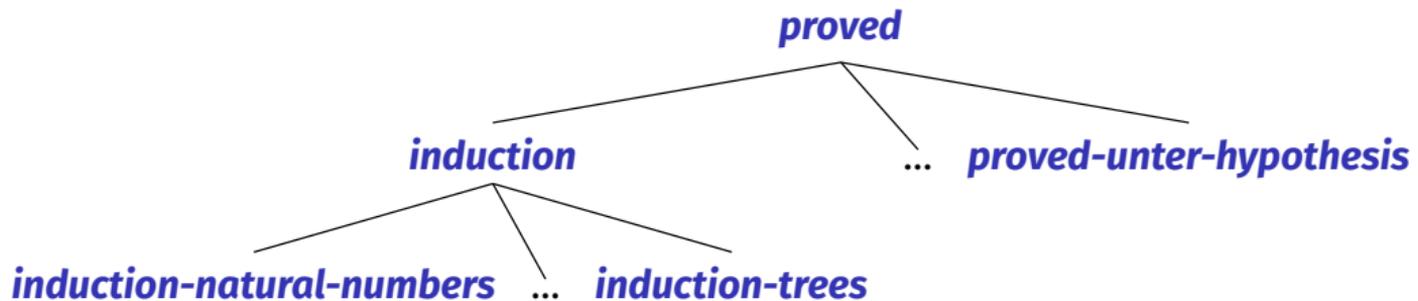
BASE-CASE

proved-under-hypothesis
 [HYPOTHESIS $\boxed{\text{h}}$]
 [THESES $\boxed{\text{t}}$]
 [ASSERTION $\boxed{\text{acc}} \Rightarrow \boxed{\text{a}}$]
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INDUCTION-STEP

proved-under-hypothesis
 [HYPOTHESIS $\boxed{\text{icond}}$: $\boxed{\text{icc}} \wedge \boxed{\text{ih}}$]
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Inheritance hierarchy of proof frames



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Further Frames

Kinds of frames and sources of information

- **Structural: Proof Techniques**, e.g. ENGEL's, 1999
- **Ontological: Domains**, e.g. MMT theories (RABE, 2016)

Another proof frame: *extremal proof*.

*"We are trying to prove the existence of an object with certain properties. The **extremal principle** tells us to pick an object which **maximizes** or **minimizes** some function. The **resulting object** is then shown to have the desired property by showing a slight perturbation (variation) would further increase or decrease the given function. [...]"*

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(ENGEL, 1999, **Problem-Solving Strategies**, p. 39)

Context and extremal proofs – interaction / ontological frames

„Das Extremalprinzip setzt also einen Kontext voraus, in dem minimale oder maximale Objekte existieren.“ (CARL, 2017, p. 75)

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Variations of extremal proofs

CARL: variation triggered by ENGEL's "three well-known facts"

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Variations of extremal proofs

CARL: variation triggered by ENGEL's "three well-known facts", e.g.

domain natural numbers: triggers *least number principle*

domain subset of reals: triggers *least upper bound principle* or *largest lower bound principle*

Context and extremal proofs – interaction by hypothesis/goal

CARL (2017, p. 75): prototypical extremal arguments are different depending on the hypothesis:

Context and extremal proofs – interaction by hypothesis/goal

CARL (2017, p. 75): prototypical extremal arguments are different depending on the hypothesis:

Beweise mithilfe des Extremalprinzips funktionieren meist auf eine der beiden folgenden Weisen (two ways):

- 1 Zu zeigen ist eine **Existenzaussage** (*existence statement*). Das extremale Objekt ist ein **Beispiel** (*example*) für ein Objekt der gesuchten Art oder hilft bei dessen **Konstruktion** (*construction*).
- 2 Zu zeigen ist eine **Allaussage** (*universal statement*). Man nimmt das **Gegenteil** (*opposite*) an, betrachtet ein extremales **Gegenbeispiel** (*counterexample*) und arbeitet auf einen **Widerspruch** (*contradiction*) (meist zur Maximalität oder Minimalität) hin.

CARL (2017, S. 75)

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Frames and Mathematical Understanding

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↪ related phenomenon: **grasping** a proof often linked to some figure of speech of **zooming out**,

↪ understanding needs knowing which frames were actually involved.

creativity: POINCARÉ saw creativity as (some fruitful) combination of old ideas or as choice among the manifold of all possible combinations.

proof identity: despite different surface structure

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- Remark: Tactics as used in proof assistants can be modeled as frames.

Outline

- 1 Introduction
- 2 Frames
- 3 Frames for Mathematical Texts
- 4 Further Frames
- 5 Frames and Mathematical Understanding
- 6 Conclusion**

Theses

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Outlook

- more frames and more texts
- corpus-based annotation workflow
- didactic experiments

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